Distributed Karger-Stein Algorithm

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The Problem: Global Minimum Cut

- Undirected, connected graph G = (V, E)
- Non-negative weights w: E -> R
- A cut is a subset S of V
- The cutsize of S is the sum of weights of edges leaving
 S
- What is the cut S of minimal cutsi



The Basic Algorithm: Karger's Randomized Algorithm

- The idea: randomly contract edges until you have a cut.
- * Probability that the result is minimal is $\Omega(1/n^2)$, so we perform $\Omega(n^2 \lg n)$ trials to get a high probability of success.
- Concretely, what do we mean by 'contract edges'?
- Give each vertex v a 'group' label g(v), and when contracting (u,v), change g(u) to g(v), and remove (u,v) from E.
- Sequential complexity: n^2 lg n trials, O(m) work per trial:
- * O(n^2m lg n)



```
g = g.mapEdges(e => e.attr.copy(rank=genRandRank()))
val max_edge = g.triplets.reduce((a, b) => max(a.rank, b.rank))
val broadcast_max_edge = sc.broadcast(max_edge)
g = g.mapVertices((_, currentGroup) => groupMapper(currentGroup, broadcast_max_edge))
        .subgraph(epred = triplet => triplet.srcAttr != triplet.dstAttr)
```



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g.edges.reduce(_ + _)

min_cut_value = 12 + 20



Data Representation: Vertex Cut





Distributed Karger: One edge contraction

Communication:

computing max rank & broadcast. Shuffle size: O(k)

Depth on each machine:

computing max rank: O(log(m/k))
assign ranks: O(1)
reassign groups: O(1)
filter edges: O(1)

Distributed Karger: n-2 edge contractions

Communication:

shuffle size: O(nk)

Depth on each machine:
 O(n log(m/k))

Distributed Karger: Total communication & Depth

Communication:

```
Total shuffle size: O(n^2 * log(n) * n * k)
```

Depth:

Total depth = depth of one trial = $O(n \log(m/k))$