Stats 300b: Theory of Statistics

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## Warning: these notes may contain factual errors

Reading: HDP Ch.8, VdV 18-19

## **Outline:**

- Sub-Gaussian Processes
- Uniform Entropy
- VC Classes

**Recap** : Process  $\{x_t\}_{t \in T}$  is p-sub-Gaussian if  $\mathbb{E}[\exp(\lambda(x_s - x_t))] \leq \frac{\lambda^2 p(s,t)^2}{2})]$  for all  $s, t \in T$ .

**Example 1:** : (Canonical symmetrized empirical process) Let  $x_i \stackrel{i.i.d}{\sim} P$  and consider  $\sup_{f \in F} (P_n f - P f)$ . Then,

$$\mathbb{E}[||P_n - P||_{\mathcal{F}}] \le 2\mathbb{E}[\sup_{f \in \mathcal{F}} |\frac{1}{n} \sum_{1}^n \epsilon_i f(x_i)] = 2\mathbb{E}[\mathbb{E}[\sup_{f \in \mathcal{F}} |\frac{1}{n} \sum_{i}^n \epsilon_i f(x_i)]|x]$$

Fix  $x_{1:n} \in \mathcal{X}^n$  and consider the process  $Z_f := \frac{1}{\sqrt{(n)}} \sum_{i=1}^n f(x_i)$ . Letting  $f, q \in \mathcal{F}$ ,

$$\mathbb{E}[exp(\lambda(Z_f - Z_g))] = \prod_{i=1}^n \mathbb{E}[\exp(\frac{\lambda}{\sqrt{n}}\epsilon_i(f(x_i) - g(x_i))] \le \exp(\frac{\lambda^2}{2n}\sum_{i=1}^n (f(x_i) - g(x_i))^2) = \exp(\lambda^2 2||f - g||_{L_1(P_n)}^2)$$

**Remark** That is,  $\{Z_f\}_{f\in\mathcal{F}}$  is a  $||\cdot||_{L_2(P_n)}$ -sub-Gaussian process. Note that  $\sup_{f\in\mathcal{F}}|\frac{1}{n}\sum_{i=1}^n$  is a sub-Gaussian process with respect to the  $L_2(P_n)$  norm, and

$$\mathbb{E}[\sup_{f \in \mathcal{F}} |P_n f - Pf|] \le \frac{1}{\sqrt{n}} 2\mathbb{E}[\mathbb{E}[\sup_{f \in \mathcal{F}} |\frac{1}{\sqrt{n}} \sum_{i=1}^n \epsilon_i f(x_i) |x|]]$$

**Goal 1** Our first goal for this lecture is to upper bound the expected suprema of sub-gaussian processes. Recall that if  $\mathcal{F}$  is bounded by B, then  $\mathbb{P}(||P_n - P||_{\mathcal{F}} \ge \mathbb{E}[||P_n - P||_{\mathcal{F}}] + 1 \le \exp(\frac{-2nt^2}{B^2})$ , using Bounded Difference, which we proved last time.

## New Material: Chaining (Dudley)

Let  $\{X_t\}_{t\in\mathcal{T}}$  be  $\rho$ -sub-Gaussian separable and mean-zero, i.e.  $\mathbb{E}[X_t] = 0$ . The idea is to control  $\sup_{t\in\mathcal{T}}X_t$  by finer and finer approximations to the supremum. We can do this because the process is separable. Let  $\mathcal{T}_0, \mathcal{T}_1, \mathcal{T}_2, \dots \mathcal{T}$  be a sequence of covers of  $\mathcal{T}$ , where  $\mathcal{T} = \min 2^{-k} \operatorname{diam}(\mathcal{T})$  cover of  $\mathcal{T}$  in the metric (or semimetric)  $\rho$ , where  $\operatorname{diam}(\mathcal{T}) := \sup_{s,t\in\mathcal{T}}\rho(s,t)$  (assumed finite),  $\mathcal{T}_0 = \{t_0\}$ , and  $\rho(t_0,t) \leq \operatorname{diam}(\mathcal{T}) \forall t \in \mathcal{T}$ .

For any  $t \in \mathcal{T}$ , consider sequences  $t_0, t_1, ..., t_k, ... \to t$  where  $t_k \in \mathcal{T}_k \ \forall k \in \mathbb{N}$ . Let  $\pi_i(t) =$ 

 $\arg\min_{t_i\in\mathcal{T}_i}\rho(t_i,t) \text{ be the closest point to } t \text{ in } \mathcal{T}_i. \text{ Fix any } k\in\mathbb{N}. \text{ Then } x_i = x_{\pi_{k-1}(t)} + x_t - x_{\pi_{k-1}(t)}.$ Let  $\pi^i(t) := \pi_i(\pi_{i+1}(\dots(\pi_{k-1}(t))\dots))$  (a concatenation of projections). Observe that

$$x_t = \sum_{i=1}^k x_{\pi_k}^i(t) - x_{\pi_k}^{i-1}(t) + x_{\pi}^0(t) = \sum_{i=1}^k x_{\pi_k}^i(t) - x_{\pi_k}^{i-1}(t) + x_{t_0}^i(t)$$

as  $\pi_k^k(t) = t$ . This is the "chain."

**Remark** For any  $k \in \mathbb{N}$ ,  $\max_{t \in \mathcal{T}} (x_t) \leq \max_{t \in \mathcal{T}} (x_{\pi_k}^i(t) - x_{\pi_k}^{i-1}(t)) + x_{\pi}^0(t)$ . How many points are there in this maximum?  $\pi_k^i(t)$  takes values in  $\mathcal{T}_i$  and  $\pi_k^{i-1}(t) = \pi_{i-1}(\pi_k^i(t))$  is a deterministic function of  $\pi_k^i(t)$ . So this is really, at "worst", a maximum over points in a set  $\mathcal{T}_i$ .

We know that if D = diam(T),  $\rho(\pi_i^i(t), \pi_k^{i-1}(t)) \le 2^{1-i}D$  as  $\pi_k^{i-1}(t) = \pi_{i-1}(\pi_k^i(t))$ ,  $T_{i-1}$  is a  $2^{1-i}$  diameter cover of T. Then,

$$\max_{t \in \mathcal{T}} x_t \le \sum_{i=1}^k \max_{t \in \mathcal{T}} (x_t - x_{\pi_{i-1}}(t)) + x_0$$

where  $t \in \mathcal{T} \max(x_t - x_{\pi_{i-1}}(t))$  is a finite maximum of  $2^{1-i}D$ -sub-Gaussian random variables. Recall that if  $\{Y_i\}_{i=1}^N$  are  $\sigma^2$ -sub-Gaussian, then

$$\mathbb{E}[\max_{i}(Y_{i})] \leq \sqrt{(2\sigma^{2}\log(N))}$$
$$\mathbb{E}[\max_{t \in T_{i}}(x_{t} - x_{\pi_{i-1}}(t))] \leq \sqrt{4^{1-i}2D^{2}\log|T_{i}|}$$

where  $Card(T_i) = \mathcal{N}(T, \rho, 2^{-i}D)$ . Then,

$$\mathbb{E}[\max_{t \in T_k}(x_t)] \le \sum_{i=1}^k \sqrt{8 \cdot 4^{-1} D^2 \log \mathcal{N}(2^{-i}D)}$$
$$= 2\sqrt{2} D \sum_{i=1}^k 2^{-i} \sqrt{\log \mathcal{N}(D, 2^{-i})}$$

Note that we can think of this as a Riemann integral, so

$$\mathbb{E}[\max_{t \in T_k}(x_t)] \le 2\sqrt{(2)}D\sum_{i=1}^k 2^{-i}\sqrt{\log \mathcal{N}(D, 2^{-i})}$$
$$\le 4\sqrt{2}D\sum_{i=1}^\infty \int_{2^{-i+1}}^{2^{-i}}\sqrt{\log \mathcal{N}(D_\epsilon)}d\epsilon$$
$$= 4\sqrt{2}D\int_0^1\sqrt{\log \mathcal{N}(D_\epsilon)}d\epsilon$$
$$= 4\sqrt{2}\int_0^{diam(T)}\sqrt{\log \mathcal{N}(T, \rho, \epsilon)}d\epsilon$$

where the last equality comes from substituting  $\epsilon$  for  $D_{\epsilon}$  and letting D = diam(T). Finally, note that  $\max_{t \in T_k \cup T_0} (x_t - x_{t_0})$  is non-negative, so Fatou's lemma implies that

$$\mathbb{E}[\sup_{t \in T_k} (x_t)] \le 4\sqrt{2} \int_0^{diam(T)} \sqrt{\log \mathcal{N}(T, \rho, \epsilon)} d\epsilon$$

**Definition 0.1.** For a metric space  $(T, \rho)$  with finite  $\rho$ -diameter  $J(T, \rho) := \int_0^{diam(T)} \sqrt{\log \mathcal{N}(T, \rho, \epsilon)} d\epsilon$  is Dudley's entropy integral.

**Theorem 1.** Let  $\{X_t\}_{t\in T}$  be a separable  $\rho$ -sub-Gaussian process. Then  $\mathbb{E}[\sup_{t\in T}(X_t)] \leq C \cdot J(T, \rho)$ , where  $C < \infty$  is a numerical constant.

**Examples** How do we control entropy integrals? (Hint: use  $\log (1 + x) \le x$ ) for small x) **Example 2:** Let  $\mathcal{F} := \{l(\theta, \cdot)\}_{\theta \in \Theta}$ , a collection of losses. For each  $x \in X$   $l(\cdot, x)$  is  $\mathcal{L}(x)$ -Lipschitz with respect to  $|| \cdot ||$  in the first argument. Assume  $\log \mathcal{N}(\Theta, || \cdot ||, \epsilon) \le d(\log (1 + \frac{diam(\Theta)}{\epsilon}))$ . We know by the entropy integral and symmetrization

$$\mathbb{E}[||P_n - P||_{\mathcal{F}}] \le C \cdot \mathbb{E}[\int_0^\infty \sqrt{\log \mathcal{N}(F, L_2(P_n), \epsilon)} d\epsilon]$$

**Remark**  $||l(t,\cdot)-l(s,\cdot)||_{L_2(P_n)} \leq \sqrt{(\log \mathcal{N}(\mathcal{F}, L_2(P_n), \epsilon)} d\epsilon \ by \ L(x) - Lipschitz.$  Thus,  $\log \mathcal{N}(\mathcal{F}, L_2(P_n), \epsilon) = 0$  if  $\epsilon \geq diam(\Theta)\sqrt{P_nL^2}$ . Also,  $\log \mathcal{N}(\mathcal{F}, ||\cdot||_{P_nL^2}, \epsilon) \leq \log \mathcal{N}(\Theta, ||\cdot||, \frac{\epsilon}{\sqrt{P_nL^2}})$ So, we have

$$\mathbb{E}[\int_{0}^{\infty} \sqrt{\log \mathcal{N}(F, L_{2}(P_{n}), \epsilon)} d\epsilon] \leq \mathbb{E}[\int_{0}^{\sqrt{P_{n}L^{2}}diam(\Theta)} \sqrt{\log \mathcal{N}(\Theta, ||\cdot||, \frac{\epsilon}{\sqrt{P_{n}L^{2}}})} d\epsilon]$$
$$= diam(\Theta)\mathbb{E}[\sqrt{P_{n}L^{2}} \int_{0}^{1} \sqrt{\log \mathcal{N}(\Theta, P_{u})} du]$$

where  $u = \frac{\epsilon}{diam\sqrt{P_nL^2}}$ 

$$\leq diam(\Theta)\mathbb{E}[L(x)^2]^{\frac{1}{2}} \int_0^1 \sqrt{d\log\left(1+\frac{1}{u}\right)} du$$
$$\leq diam(\Theta)\mathbb{E}[L(x)^2]^{\frac{1}{2}} \int_0^1 \sqrt{\frac{d}{u}} du$$
$$\leq diam(\Theta)\mathbb{E}[L(x)^2]^{\frac{1}{2}} \sqrt{d}$$

**Next Goal** Give classes  $\mathcal{F}$  for which we can bound  $\sup_{Q} \mathcal{N}(\mathcal{F}, L_2(Q), \epsilon)$ .

VC Classes Big example of classes allowing uniform bounds on entropy numbers.

**Definition 0.2.** Let C be a collection of sets and  $X = x_1, ..., x_n$ . A vector  $y \in \{\pm 1\}^n$  is a labeling of X. We say C shatters X if for all labelings  $y \in \{\pm 1\}^n$ ,  $\exists a \text{ set } A \in C$  such that  $x_i \in A$  if  $y_i = +1$  and  $x_i \notin A$  if  $y_i = -1$ .

**Example 3:** Let  $x_1, x_2, x_3 \in \mathbb{R}^3$  not collinear. C=Half-spaces in  $\mathbb{R}^2$ . For any labeling, these points can be shattered.

**Definition 0.3.** Given  $\mathcal{C} \subset 2^{\mathcal{X}}$ , the shattering coefficient of  $\mathcal{C}$  on  $x_1, x_2, ..., x_n$  is  $\Delta_n(\mathcal{C}, x_{1:n}) := card\{A \cap x_1, ..., x_n : A \in \mathcal{C}\} = the number of labelings of <math>x_{1:n}$  that  $\mathcal{C}$  gives.

The VC-dimension (Vapnik-Chervonenkis) of C is  $VC(C) := \sup\{n \in \mathbb{N} : \max_{X_{1:n} \in \mathcal{X}^n} \Delta_n(C, x_{1:n}) = 2^n\} = \text{the size of the largest set of points that } C$ ca shatter.

**Lemma 2.** Sauer-Shelah lemma For any class C of sets,

$$\max_{x_{1:n}\in X^n} \Delta_n(\mathcal{C}, x_{1:n}) \le \sum_{j=0}^{VC(\mathcal{C})} \binom{n}{j} = O(n^{VC(\mathcal{C})})$$

**Consequence:** If  $\max_{x_{1:n} \in X^n} \Delta_n(\mathcal{C}, x_{1:n}) < 2^n$ , then  $VC(\mathcal{C}) < n$  and

$$\Delta_n(\mathcal{C}, x_{1:n}) \le O(1) \cdot n^{VC(\mathcal{C})}$$

. Additional lectures notes on the course website provide a further reference on this topic.