

Lecture 2

LQR via Lagrange multipliers

- useful matrix identities
- linearly constrained optimization
- LQR via constrained optimization

Some useful matrix identities

let's start with a simple one:

$$Z(I + Z)^{-1} = I - (I + Z)^{-1}$$

(provided $I + Z$ is invertible)

to verify this identity, we start with

$$I = (I + Z)(I + Z)^{-1} = (I + Z)^{-1} + Z(I + Z)^{-1}$$

re-arrange terms to get identity

an identity that's a bit more complicated:

$$(I + XY)^{-1} = I - X(I + YX)^{-1}Y$$

(if either inverse exists, then the other does; in fact
 $\det(I + XY) = \det(I + YX)$)

to verify:

$$\begin{aligned}(I - X(I + YX)^{-1}Y) (I + XY) &= I + XY - X(I + YX)^{-1}Y(I + XY) \\ &= I + XY - X(I + YX)^{-1}(I + YX)Y \\ &= I + XY - XY = I\end{aligned}$$

another identity:

$$Y(I + XY)^{-1} = (I + YX)^{-1}Y$$

to verify this one, start with $Y(I + XY) = (I + YX)Y$

then multiply on left by $(I + YX)^{-1}$, on right by $(I + XY)^{-1}$

- note dimensions of inverses not necessarily the same
- mnemonic: lefthand Y moves into inverse, pushes righthand Y out . . .

and one more:

$$(I + XZ^{-1}Y)^{-1} = I - X(Z + YX)^{-1}Y$$

let's check:

$$\begin{aligned}(I + X(Z^{-1}Y))^{-1} &= I - X(I + Z^{-1}YX)^{-1}Z^{-1}Y \\ &= I - X(Z(I + Z^{-1}YX))^{-1}Y \\ &= I - X(Z + YX)^{-1}Y\end{aligned}$$

Example: rank one update

- suppose we've already calculated or know A^{-1} , where $A \in \mathbf{R}^{n \times n}$
- we need to calculate $(A + bc^T)^{-1}$, where $b, c \in \mathbf{R}^n$
($A + bc^T$ is called a *rank one update* of A)

we'll use another identity, called *matrix inversion lemma*:

$$(A + bc^T)^{-1} = A^{-1} - \frac{1}{1 + c^T A^{-1} b} (A^{-1} b)(c^T A^{-1})$$

note that RHS is easy to calculate since we know A^{-1}

more general form of matrix inversion lemma:

$$(A + BC)^{-1} = A^{-1} - A^{-1}B(I + CA^{-1}B)^{-1}CA^{-1}$$

let's verify it:

$$\begin{aligned}(A + BC)^{-1} &= (A(I + A^{-1}BC))^{-1} \\ &= (I + (A^{-1}B)C)^{-1}A^{-1} \\ &= (I - (A^{-1}B)(I + C(A^{-1}B))^{-1}C) A^{-1} \\ &= A^{-1} - A^{-1}B(I + CA^{-1}B)^{-1}CA^{-1}\end{aligned}$$

Another formula for the Riccati recursion

$$\begin{aligned}P_{t-1} &= Q + A^T P_t A - A^T P_t B (R + B^T P_t B)^{-1} B^T P_t A \\&= Q + A^T P_t (I - B (R + B^T P_t B)^{-1} B^T P_t) A \\&= Q + A^T P_t (I - B ((I + B^T P_t B R^{-1}) R)^{-1} B^T P_t) A \\&= Q + A^T P_t (I - B R^{-1} (I + B^T P_t B R^{-1})^{-1} B^T P_t) A \\&= Q + A^T P_t (I + B R^{-1} B^T P_t)^{-1} A \\&= Q + A^T (I + P_t B R^{-1} B^T)^{-1} P_t A\end{aligned}$$

or, in pretty, symmetric form:

$$P_{t-1} = Q + A^T P_t^{1/2} \left(I + P_t^{1/2} B R^{-1} B^T P_t^{1/2} \right)^{-1} P_t^{1/2} A$$

Linearly constrained optimization

$$\begin{array}{ll} \text{minimize} & f(x) \\ \text{subject to} & Fx = g \end{array}$$

- $f : \mathbf{R}^n \rightarrow \mathbf{R}$ is smooth *objective function*
- $F \in \mathbf{R}^{m \times n}$ is fat

form *Lagrangian* $L(x, \lambda) = f(x) + \lambda^T(g - Fx)$ (λ is *Lagrange multiplier*)

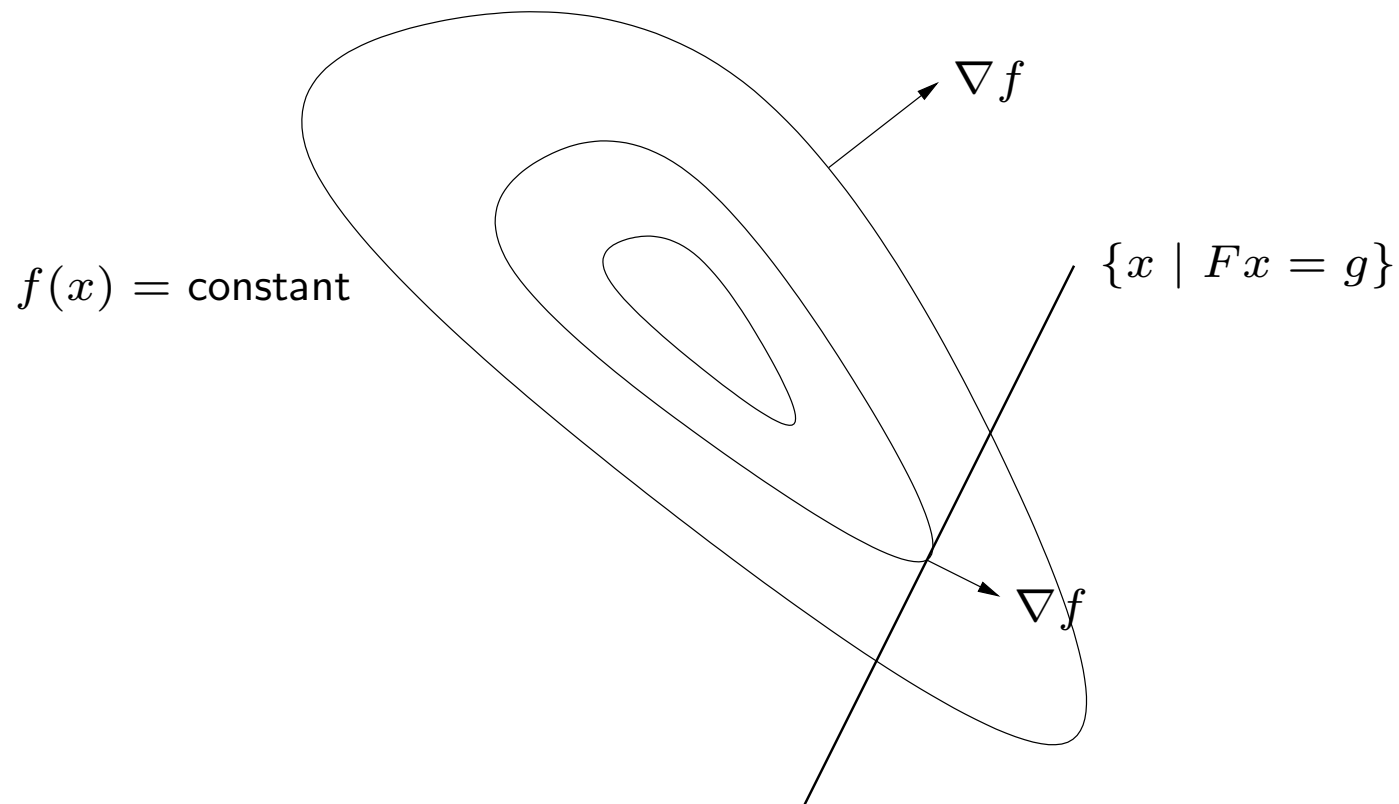
if x is optimal, then

$$\nabla_x L = \nabla f(x) - F^T \lambda = 0, \quad \nabla_\lambda L = g - Fx = 0$$

i.e., $\nabla f(x) = F^T \lambda$ for some $\lambda \in \mathbf{R}^m$

(generalizes optimality condition $\nabla f(x) = 0$ for unconstrained minimization problem)

Picture



$$\nabla f(x) = F^T \lambda \text{ for some } \lambda \iff \nabla f(x) \in \mathcal{R}(F^T) \iff \nabla f(x) \perp \mathcal{N}(F)$$

Feasible descent direction

suppose x is current, feasible point (*i.e.*, $Fx = g$)

consider a small step in direction v , to $x + hv$ (h small, positive)

when is $x + hv$ better than x ?

need $x + hv$ feasible: $F(x + hv) = g + hFv = g$, so $Fv = 0$

$v \in \mathcal{N}(F)$ is called a *feasible direction*

we need $x + hv$ to have smaller objective than x :

$$f(x + hv) \approx f(x) + h\nabla f(x)^T v < f(x)$$

so we need $\nabla f(x)^T v < 0$ (called a *descent direction*)

(if $\nabla f(x)^T v > 0$, $-v$ is a descent direction, so we need only $\nabla f(x)^T v \neq 0$)

x is not optimal if there exists a feasible descent direction

if x is optimal, every feasible direction satisfies $\nabla f(x)^T v = 0$

$$\begin{aligned} Fv = 0 \Rightarrow \nabla f(x)^T v = 0 &\iff \mathcal{N}(F) \subseteq \mathcal{N}(\nabla f(x)^T) \\ &\iff \mathcal{R}(F^T) \supseteq \mathcal{R}(\nabla f(x)) \\ &\iff \nabla f(x) \in \mathcal{R}(F^T) \\ &\iff \nabla f(x) = F^T \lambda \text{ for some } \lambda \in \mathbf{R}^m \\ &\iff \nabla f(x) \perp \mathcal{N}(F) \end{aligned}$$

LQR as constrained minimization problem

$$\begin{aligned} \text{minimize} \quad & J = \frac{1}{2} \sum_{t=0}^{N-1} (x_t^T Q x_t + u_t^T R u_t) + \frac{1}{2} x_N^T Q_f x_N \\ \text{subject to} \quad & x_{t+1} = A x_t + B u_t, \quad t = 0, \dots, N-1 \end{aligned}$$

- variables are u_0, \dots, u_{N-1} and x_1, \dots, x_N
($x_0 = x^{\text{init}}$ is given)
- objective is (convex) quadratic
(factor $1/2$ in objective is for convenience)

introduce Lagrange multipliers $\lambda_1, \dots, \lambda_N \in \mathbf{R}^n$ and form Lagrangian

$$L = J + \sum_{t=0}^{N-1} \lambda_{t+1}^T (A x_t + B u_t - x_{t+1})$$

Optimality conditions

we have $x_{t+1} = Ax_t + Bu_t$ for $t = 0, \dots, N - 1$, $x_0 = x^{\text{init}}$

for $t = 0, \dots, N - 1$, $\nabla_{u_t} L = Ru_t + B^T \lambda_{t+1} = 0$

hence, $u_t = -R^{-1} B^T \lambda_{t+1}$

for $t = 1, \dots, N - 1$, $\nabla_{x_t} L = Qx_t + A^T \lambda_{t+1} - \lambda_t = 0$

hence, $\lambda_t = A^T \lambda_{t+1} + Qx_t$

$\nabla_{x_N} L = Q_f x_N - \lambda_N = 0$, so $\lambda_N = Q_f x_N$

these are a set of linear equations in the variables

$$u_0, \dots, u_{N-1}, \quad x_1, \dots, x_N, \quad \lambda_1, \dots, \lambda_N$$

Co-state equations

optimality conditions break into two parts:

$$x_{t+1} = Ax_t + Bu_t, \quad x_0 = x^{\text{init}}$$

this recursion for state x runs forward in time, with initial condition

$$\lambda_t = A^T \lambda_{t+1} + Qx_t, \quad \lambda_N = Q_f x_N$$

this recursion for λ runs backward in time, with final condition

- λ is called *co-state*
- recursion for λ sometimes called *adjoint system*

Solution via Riccati recursion

we will see that $\lambda_t = P_t x_t$, where P_t is the min-cost-to-go matrix defined by the Riccati recursion

thus, Riccati recursion gives clever way to solve this set of linear equations

it holds for $t = N$, since $P_N = Q_f$ and $\lambda_N = Q_f x_N$

now suppose it holds for $t + 1$, *i.e.*, $\lambda_{t+1} = P_{t+1} x_{t+1}$

let's show it holds for t , *i.e.*, $\lambda_t = P_t x_t$

using $x_{t+1} = Ax_t + Bu_t$ and $u_t = -R^{-1} B^T \lambda_{t+1}$,

$$\lambda_{t+1} = P_{t+1}(Ax_t + Bu_t) = P_{t+1}(Ax_t - BR^{-1}B^T \lambda_{t+1})$$

so

$$\lambda_{t+1} = (I + P_{t+1}BR^{-1}B^T)^{-1}P_{t+1}Ax_t$$

using $\lambda_t = A^T \lambda_{t+1} + Qx_t$, we get

$$\lambda_t = A^T (I + P_{t+1} B R^{-1} B^T)^{-1} P_{t+1} A x_t + Q x_t = P_t x_t$$

since by the Riccati recursion

$$P_t = Q + A^T (I + P_{t+1} B R^{-1} B^T)^{-1} P_{t+1} A$$

this proves $\lambda_t = P_t x_t$

let's check that our two formulas for u_t are consistent:

$$\begin{aligned}u_t &= -R^{-1}B^T\lambda_{t+1} \\ &= -R^{-1}B^T(I + P_{t+1}BR^{-1}B^T)^{-1}P_{t+1}Ax_t \\ &= -R^{-1}(I + B^TP_{t+1}BR^{-1})^{-1}B^TP_{t+1}Ax_t \\ &= -((I + B^TP_{t+1}BR^{-1})R)^{-1}B^TP_{t+1}Ax_t \\ &= -(R + B^TP_{t+1}B)^{-1}B^TP_{t+1}Ax_t\end{aligned}$$

which is what we had before