# Balanced Label Propagation for Partitioning Massive Graphs 

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## Goal: partition a really big graph


facebook

## Motivation: distributed computation

- Distributing graph calculations ('sharding a graph’) makes traversal/aggregation very expensive.



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Oh... and the algorithm better be FAST.

## Partitioning a really big graph: How?

- Garey, Johnson, Stockmeyer 1976: Minimum bisection is NP-hard
- Karypsis and Kumar 1998:

METIS

- Feige and Krautgamer 2000:
$\mathrm{O}\left(\mathrm{n}^{1 / 2} \log \mathrm{n}\right)$-factor approximation


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$\mathrm{O}\left(\mathrm{n}^{1 / 2} \log \mathrm{n}\right)$-factor approximation
- METIS does not scale to 100B+ edges.
- Need a principled approach, ideally one that can be Hadoop-ified.


## Basic idea: Label propagation

- Iteratively move nodes to be with the plurality of their neighbors:



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## Balance via Linear Program

- Greedily maximize edge locality with constraints (max/min sizes $\mathrm{s}_{\mathrm{i}} \mathrm{T}_{\mathrm{i}}$ ):
$x_{i j} \quad$ Solution: number of people to move from i to j .
$f_{i j}(x) \quad$ Cumulative gain from moving $x$ people (ordered by co-location gain).



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\max _{X} \sum_{i, j} f_{i j}\left(x_{i j}\right) \quad \text { s.t. }\left\{\begin{array}{rcccc}
S_{i}-\left|V_{i}\right| & \leq & \sum_{j \neq i}\left(x_{i j}-x_{j i}\right) & \leq & T_{i}-\left|V_{i}\right|, \\
0 & \leq & x_{i j} & \leq & P_{i j}, \\
& \forall i, j
\end{array}\right.
$$

(Maximize the co-location gain of all machine swaps)
(Subject to balance)
(and the number of people available to move)

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- Linear Program: n=78 machines => 12k variables / 400k constraints

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\max _{X, Z} \sum_{i, j} z_{i j} \quad \text { s.t. }\left\{\begin{array}{rrrr}
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x_{i j} & \leq i \\
0 \leq & P_{i j}, & \forall i, j \\
-a_{i j k} x_{i j}+z_{i j} & \leq & b_{i j k}, & \forall i, j, k
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- Summary of algorithm:
- Step 1: Figure out who wants to move
- Step 2: Solve LP to decide who can move without breaking balance
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Step 2 is the contribution compared to ordinary Label Prop.

## What about geography?


facebook

## Initialization using geography

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- Possible to do much better than random with Facebook, using geography.
- Spatial model of small-world networks (for routing): Kleinberg 2000
- Validation: Liben-Nowell et al. 2005; Backstrom, Sun, Marlow 2010.
- Friendship probability as a function of rank-distance:



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- Geographic initialization 'converges' within 1 step
- Random initialization slow to start when: avg degree >\# partitions Use 'restraint': only move big gainers (*s below)
- LJ partitioning quality not so dependent on \# partitions:

BLP exploiting primarily local structure.



## Results: Machine adjacency matrix



Random Initialization + Prop


Geographic Initialization


- Random initialization + 8 step prop
- Geographic initialization ONLY
- Geographic + 1 step prop
- Targeting n=78 machines:

2 racks of 39 , visible as blocks

## ‘People You May Know’

- PYMK = ‘People You May Know’
- Ranked suggestion of friends-of-friends (FoFs) as friends.
- Average user has 40 k FoFs, widely distributed.
- Ranks 145,000,000 suggestions per second.
- Graph distributed across 78 machines with 72GB RAM each.


Friends of Friends

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- Query time: What about overhead? Faster or slower?


## Results: Query time / network traffic

- Median number of machines hit per query reduced from 60 to 9.
- Query time reduced by 49\%, traffic reduced by 63\%:




## Conclusions and Future work

- Label propagation is fast, we show it can be constrained
- Social networks very clustered, making local algorithms very effective
- Geographic metadata very useful
- Sharding greatly improves distributed graph computations such as PYMK

