Balanced Label Propagation for Partitioning Massive Graphs

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Goal: partition a really big graph



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Oh... and the algorithm better be FAST.

Partitioning a really big graph: How?

- Garey, Johnson, Stockmeyer 1976:
- Karypsis and Kumar 1998:
- Feige and Krautgamer 2000:

Minimum bisection is NP-hard METIS O(n^{1/2} log n)-factor approximation

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- METIS does not scale to 100B+ edges.
- Need a principled approach, ideally one that can be Hadoop-ified.

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- Greedily maximize edge locality with constraints (max/min sizes S_i,T_i):

- x_{ij} Solution: number of people to move from i to j.
- $f_{ij}(x)$ Cumulative gain from moving x people (ordered by co-location gain).



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$$\max_{X} \sum_{i,j} f_{ij}(x_{ij}) \qquad \text{s.t.} \begin{cases} S_i - |V_i| \leq \sum_{j \neq i} (x_{ij} - x_{ji}) \leq T_i - |V_i|, & \forall i \\ 0 \leq x_{ij} \leq P_{ij}, & \forall i, j \end{cases}$$

(Maximize the co-location gain of all machine swaps)

(Subject to balance) (and the number of people available to move)

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Linear Program: n=78 machines => 12k variables / 400k constraints

$$\max_{X,Z} \sum_{i,j} z_{ij} \qquad \text{s.t.} \begin{cases} S_i - |V_i| \leq \sum_{j \neq i} (x_{ij} - x_{ji}) \leq T_i - |V_i|, & \forall i \\ 0 \leq x_{ij} \leq P_{ij}, & \forall i, j \\ -a_{ijk} x_{ij} + z_{ij} \leq b_{ijk}, & \forall i, j, k \end{cases}$$

Summary of algorithm:

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- Step 2: Solve LP to decide who can move without breaking balance
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Step 2 is the contribution compared to ordinary Label Prop.

What about geography?



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 - Spatial model of small-world networks (for routing): Kleinberg 2000
 - Validation: Liben-Nowell et al. 2005; Backstrom, Sun, Marlow 2010.
 - Friendship probability as a function of rank-distance:



- Backstrom, Sun, Marlow 2010

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- Geographic initialization 'converges' within 1 step
- Random initialization slow to start when: avg degree > # partitions
 Use 'restraint': only move big gainers (*s below)
- LJ partitioning quality not so dependent on # partitions: BLP exploiting primarily local structure.



Results: Machine adjacency matrix



- Random initialization + 8 step prop
- Geographic initialization ONLY
- Geographic + 1 step prop
- Targeting n=78 machines:
 2 racks of 39, visible as blocks



'People You May Know'

- PYMK = 'People You May Know'
- Ranked suggestion of friends-of-friends (FoFs) as friends.
- Average user has 40k FoFs, widely distributed.
- Ranks 145,000,000 suggestions per second.
- Graph distributed across 78 machines with 72GB RAM each.



Friends of Friends

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Friends of Friends

Me = ID #3

Friends IDs: 24, 54, 4, 30, 8

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• Query time: What about overhead? Faster or slower?

Results: Query time / network traffic

- Median number of machines hit per query reduced from 60 to 9.
- Query time reduced by 49%, traffic reduced by 63%:



Conclusions and Future work

- Label propagation is fast, we show it can be constrained
- Social networks very clustered, making local algorithms very effective
- Geographic metadata very useful

Sharding greatly improves distributed graph computations such as PYMK