

Ramsey numbers of link hypergraphs

Xiaoyu He
Stanford University

Joint work with Jacob Fox

December 12, 2019

Definition

A *k*-uniform hypergraph, or simply *k*-graph, $H = (V, E)$ is a set V of vertices and a set $E \subseteq \binom{V}{k}$ of edges.

Hypergraphs

Definition

A k -uniform hypergraph, or simply k -graph, $H = (V, E)$ is a set V of vertices and a set $E \subseteq \binom{V}{k}$ of edges.

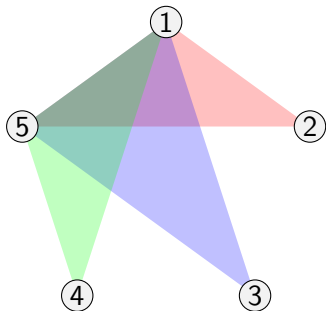


Figure: A 3-graph on 5 vertices and 3 edges.

Hypergraph Ramsey numbers

Definition

The *Ramsey number* $r(H, G)$ of two k -graphs H and G is the smallest N such that for any k -graph Γ on N vertices, either $H \subset \Gamma$ or $G \subset \bar{\Gamma}$.

Hypergraph Ramsey numbers

Definition

The *Ramsey number* $r(H, G)$ of two k -graphs H and G is the smallest N such that for any k -graph Γ on N vertices, either $H \subset \Gamma$ or $G \subset \bar{\Gamma}$.

For example, $r(K_3, K_3) = 6$.

Hypergraph Ramsey numbers

Definition

The *Ramsey number* $r(H, G)$ of two k -graphs H and G is the smallest N such that for any k -graph Γ on N vertices, either $H \subset \Gamma$ or $G \subset \bar{\Gamma}$.

For example, $r(K_3, K_3) = 6$.

Theorem (Ramsey 1930)

For any $k \geq 1$ and k -graphs H and G , $r(H, G) < \infty$.

Hypergraph Ramsey numbers

Definition

The *Ramsey number* $r(H, G)$ of two k -graphs H and G is the smallest N such that for any k -graph Γ on N vertices, either $H \subset \Gamma$ or $G \subset \bar{\Gamma}$.

For example, $r(K_3, K_3) = 6$.

Theorem (Ramsey 1930)

For any $k \geq 1$ and k -graphs H and G , $r(H, G) < \infty$.

Theorem (Erdős-Szekeres 1935)

In any sufficiently large set of points in general position in the plane, some n form a convex polygon.

Hypergraph Ramsey numbers

Definition

The *Ramsey number* $r(H, G)$ of two k -graphs H and G is the smallest N such that for any k -graph Γ on N vertices, either $H \subset \Gamma$ or $G \subset \bar{\Gamma}$.

For example, $r(K_3, K_3) = 6$.

Theorem (Ramsey 1930)

For any $k \geq 1$ and k -graphs H and G , $r(H, G) < \infty$.

Theorem (Erdős-Szekeres 1935)

In any sufficiently large set of points in general position in the plane, some n form a convex polygon.

Remark: this can be deduced from $r(K_n^{(3)}, K_n^{(3)}) < \infty$ or from $r(K_5^{(4)}, K_n^{(4)}) < \infty$.

- 1. Most hypergraph Ramsey problems reduce to uniformity 3, but not uniformity 2.**

- 1. Most hypergraph Ramsey problems reduce to uniformity 3, but not uniformity 2.**
- 2. Quasirandomness conditions for hypergraphs are not all equivalent.**

- 1. Most hypergraph Ramsey problems reduce to uniformity 3, but not uniformity 2.**
- 2. Quasirandomness conditions for hypergraphs are not all equivalent.**
- 3. We can design Ramsey hypergraphs that are globally quasirandom but locally structured.**

Diagonal Ramsey numbers

Theorem (Erdős 1947, Erdős-Szekeres 1935)

$$2^{n/2} \leq r(K_n, K_n) \leq 2^{2n}.$$

Diagonal Ramsey numbers

Theorem (Erdős 1947, Erdős-Szekeres 1935)

$$2^{n/2} \leq r(K_n, K_n) \leq 2^{2n}.$$

The lower bound construction is the random graph.

Diagonal Ramsey numbers

Theorem (Erdős 1947, Erdős-Szekeres 1935)

$$2^{n/2} \leq r(K_n, K_n) \leq 2^{2n}.$$

The lower bound construction is the random graph.

Theorem (Erdős-Hajnal 1960s, Erdős-Rado 1952)

For $k \geq 3$,

$$t_{k-1}(\Omega(n^2)) \leq r(K_n^{(k)}, K_n^{(k)}) \leq t_k(O(n)),$$

where $t_k(n)$ is the tower function $t_1(n) = n$, $t_{k+1}(n) = 2^{t_k(n)}$.

Diagonal Ramsey numbers

Theorem (Erdős 1947, Erdős-Szekeres 1935)

$$2^{n/2} \leq r(K_n, K_n) \leq 2^{2n}.$$

The lower bound construction is the random graph.

Theorem (Erdős-Hajnal 1960s, Erdős-Rado 1952)

For $k \geq 3$,

$$t_{k-1}(\Omega(n^2)) \leq r(K_n^{(k)}, K_n^{(k)}) \leq t_k(O(n)),$$

where $t_k(n)$ is the tower function $t_1(n) = n$, $t_{k+1}(n) = 2^{t_k(n)}$.

Remark

Both the upper and lower bounds are recursive in nature, proving bounds on uniformity $k + 1$ using uniformity k . However, the lower bound (stepping up lemma) only works starting from $k = 3$.

Off-diagonal Ramsey numbers

In the graph case:

Theorem (Kim 1995, Ajtai-Komlós-Szemerédi 1980)

$$r(K_3, K_n) = \Theta\left(\frac{n^2}{\log n}\right).$$

In the graph case:

Theorem (Kim 1995, Ajtai-Komlós-Szemerédi 1980)

$$r(K_3, K_n) = \Theta\left(\frac{n^2}{\log n}\right).$$

Central problem in the development of the probabilistic method:

- Alterations (Erdős 1961)
- Lovász Local Lemma (Spencer 1975)
- Large deviation inequalities (Krivelevich 1995)
- Rödl nibble (Kim 1995)
- The H -free process (Erdős-Suen-Winkler 1995, Bohman-Keevash 2010)

In the graph case:

Theorem (Kim 1995, Ajtai-Komlós-Szemerédi 1980)

$$r(K_3, K_n) = \Theta\left(\frac{n^2}{\log n}\right).$$

Central problem in the development of the probabilistic method:

- Alterations (Erdős 1961)
- Lovász Local Lemma (Spencer 1975)
- Large deviation inequalities (Krivelevich 1995)
- Rödl nibble (Kim 1995)
- The H -free process (Erdős-Suen-Winkler 1995, Bohman-Keevash 2010)

For almost all other H , the order of $r(H, K_n)$ is still unknown.

Off-diagonal hypergraph Ramsey numbers

Let $K_4^{(3)} - e$ be the 3-graph with 4 vertices and 3 edges.

Theorem (Erdős-Hajnal 1972)

$$2^{\Omega(n)} \leq r(K_4^{(3)} - e, K_n^{(3)}) \leq 2^{O(n \log n)}.$$

Off-diagonal hypergraph Ramsey numbers

Let $K_4^{(3)} - e$ be the 3-graph with 4 vertices and 3 edges.

Theorem (Erdős-Hajnal 1972)

$$2^{\Omega(n)} \leq r(K_4^{(3)} - e, K_n^{(3)}) \leq 2^{O(n \log n)}.$$

This was also the best known lower bound for $r(K_4^{(3)}, K_n^{(3)})$ until:

Theorem (Conlon-Fox-Sudakov 2010)

$$2^{\Omega(n \log n)} \leq r(K_4^{(3)}, K_n^{(3)}) \leq 2^{O(n^2 \log n)}.$$

Off-diagonal hypergraph Ramsey numbers

Let $K_4^{(3)} - e$ be the 3-graph with 4 vertices and 3 edges.

Theorem (Erdős-Hajnal 1972)

$$2^{\Omega(n)} \leq r(K_4^{(3)} - e, K_n^{(3)}) \leq 2^{O(n \log n)}.$$

This was also the best known lower bound for $r(K_4^{(3)}, K_n^{(3)})$ until:

Theorem (Conlon-Fox-Sudakov 2010)

$$2^{\Omega(n \log n)} \leq r(K_4^{(3)}, K_n^{(3)}) \leq 2^{O(n^2 \log n)}.$$

Our main result:

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Remark

A purely random 3-graph on N vertices does poorly, since edge density $p = N^{-c}$, which makes the independence number $\approx N^{c/2}$.

The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament
on $N = 2^{cn}$ vertices:

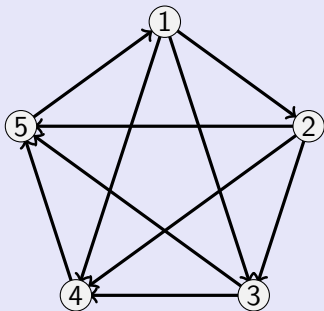
The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament on $N = 2^{cn}$ vertices:



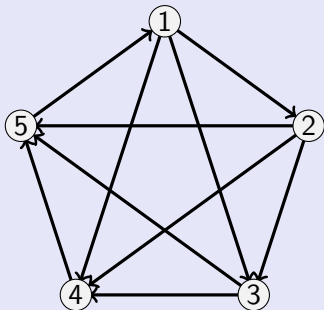
The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

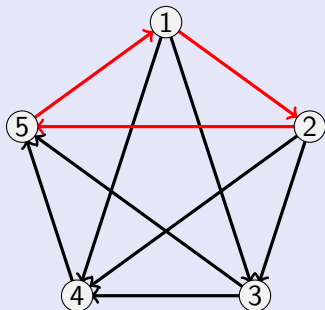
$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament on $N = 2^{cn}$ vertices:



Let Γ be the 3-graph of cyclic triangles in T :



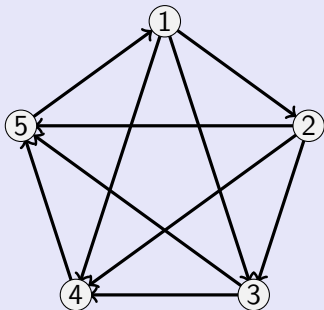
The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

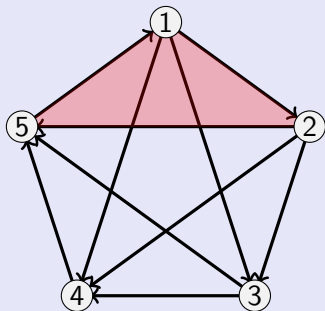
$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament on $N = 2^{cn}$ vertices:



Let Γ be the 3-graph of cyclic triangles in T :



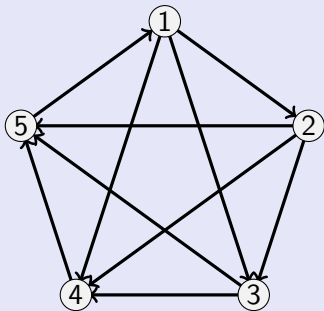
The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

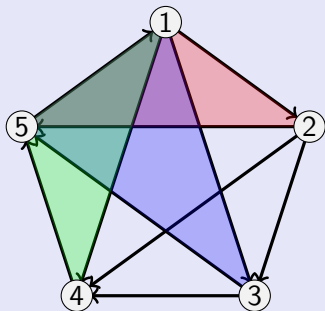
$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament on $N = 2^{cn}$ vertices:



Let Γ be the 3-graph of cyclic triangles in T :



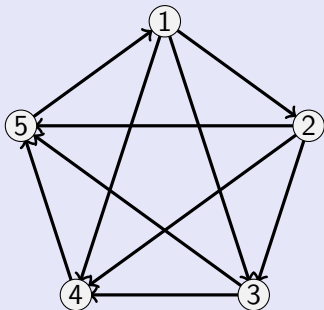
The Erdős-Hajnal tournament construction

Theorem (Erdős-Hajnal 1972)

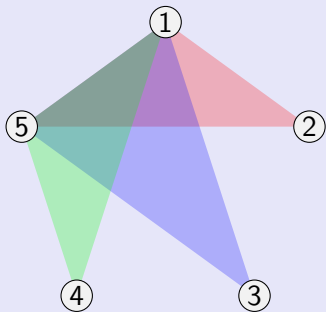
$$r(K_4^{(3)} - e, K_n^{(3)}) \geq 2^{\Omega(n)}.$$

Proof.

Let T be a random tournament on $N = 2^{cn}$ vertices:



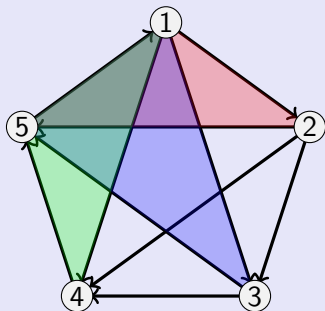
Let Γ be the 3-graph of cyclic triangles in T :



The Erdős-Hajnal tournament construction

Proof (continued).

Among any four vertices in T , at most two out of four triples form cyclic triangles, so Γ doesn't contain $K_4^{(3)} - e$.



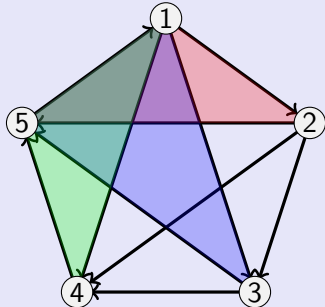
The Erdős-Hajnal tournament construction

Proof (continued).

Among any four vertices in T , at most two out of four triples form cyclic triangles, so Γ doesn't contain $K_4^{(3)} - e$.

The expected number of independent sets of size n in Γ is

$$\binom{N}{n} \cdot n! \cdot 2^{-\binom{n}{2}} < 1.$$



Theorem (Chung, Graham, Wilson 1989)

If H is a fixed labelled graph, G is a labelled graph on n vertices, and every vertex subset $U \subseteq V(G)$ contains $p \binom{|U|}{2} + o(n^2)$ edges, then G contains $(1 + o(1))p^{e(H)} n^{v(H)}$ labelled copies of H .

Theorem (Chung, Graham, Wilson 1989)

If H is a fixed labelled graph, G is a labelled graph on n vertices, and every vertex subset $U \subseteq V(G)$ contains $p \binom{|U|}{2} + o(n^2)$ edges, then G contains $(1 + o(1))p^{e(H)} n^{v(H)}$ labelled copies of H .

Example: counting triangles

If every linear-sized subset of a graph G has edge density $1/4$, then G has triangle density $1/64$.

Digression: Quasirandomness for hypergraphs

Theorem (Chung, Graham, Wilson 1989)

If H is a fixed labelled graph, G is a labelled graph on n vertices, and every vertex subset $U \subseteq V(G)$ contains $p \binom{|U|}{2} + o(n^2)$ edges, then G contains $(1 + o(1))p^{e(H)} n^{v(H)}$ labelled copies of H .

Example: counting triangles

If every linear-sized subset of a graph G has edge density $1/4$, then G has triangle density $1/64$.

Surprise

Such a statement is false for 3-graphs!

Digression: Quasirandomness for hypergraphs

If Γ is the 3-graph of cyclic triangles in a random tournament T on N vertices, then every subset $U \subseteq V(G)$ contains

$$\frac{1}{4} \binom{|U|}{3} + o(N^3)$$

edges, and yet Γ is $K_4^{(3)}$ -free.

Digression: Quasirandomness for hypergraphs

If Γ is the 3-graph of cyclic triangles in a random tournament T on N vertices, then every subset $U \subseteq V(G)$ contains

$$\frac{1}{4} \binom{|U|}{3} + o(N^3)$$

edges, and yet Γ is $K_4^{(3)}$ -free.

Remarks

- This is a type of random construction that isn't available in uniformity 2: it is quasirandom in the sense of edge densities and non-quasirandom in the sense of subgraph counts.

Digression: Quasirandomness for hypergraphs

If Γ is the 3-graph of cyclic triangles in a random tournament T on N vertices, then every subset $U \subseteq V(G)$ contains

$$\frac{1}{4} \binom{|U|}{3} + o(N^3)$$

edges, and yet Γ is $K_4^{(3)}$ -free.

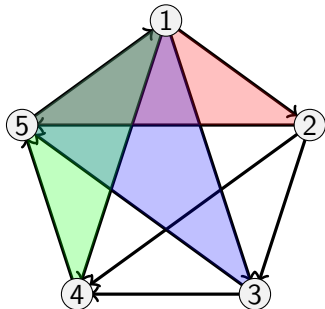
Remarks

- This is a type of random construction that isn't available in uniformity 2: it is quasirandom in the sense of edge densities and non-quasirandom in the sense of subgraph counts.
- The existence of such hypergraphs has serious implications for hypergraph regularity.

Improving the tournament lower bound

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$



Definition

If G is a k -graph, the *link* G_v of a vertex v in G is the $(k - 1)$ -graph on $V(G) \setminus \{v\}$ whose edges come from deleting v from the edges of G containing v .

Links in hypergraphs

Definition

If G is a k -graph, the *link* G_v of a vertex v in G is the $(k-1)$ -graph on $V(G) \setminus \{v\}$ whose edges come from deleting v from the edges of G containing v .

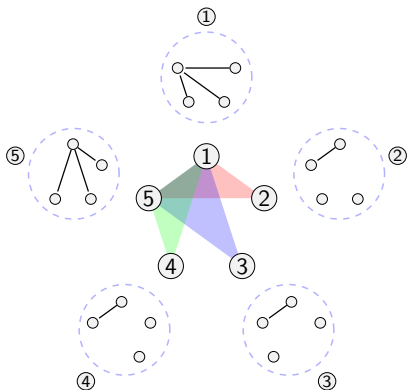


Figure: The links of each vertex.

Links in hypergraphs

Definition

If G is a k -graph, the *link* G_v of a vertex v in G is the $(k-1)$ -graph on $V(G) \setminus \{v\}$ whose edges come from deleting v from the edges of G containing v .

Note: G is $(K_4^{(3)} - e)$ -free iff the links G_v are all triangle-free.

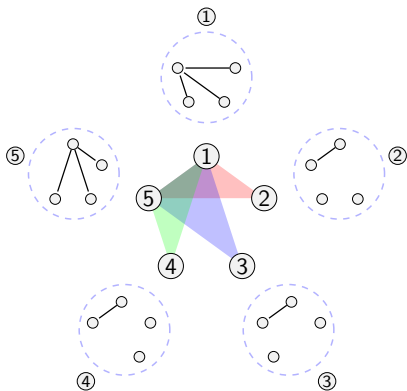


Figure: The links of each vertex.

Links of the tournament construction

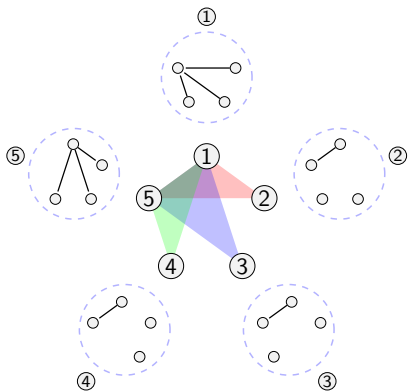


Figure: The links of each vertex.

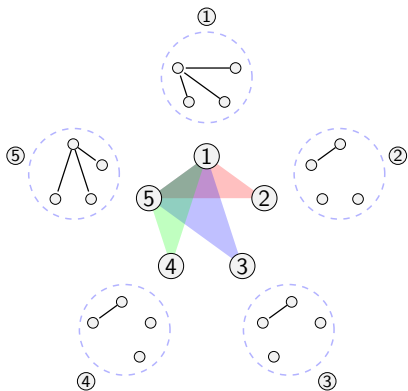


Figure: The links of each vertex.

Observations

- 1 If Γ is the 3-graph of cyclic triangles in any tournament, then the links Γ_v of its vertices are bipartite.

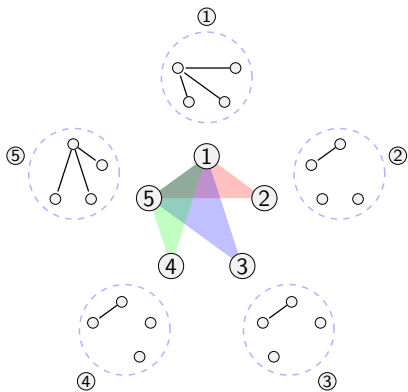


Figure: The links of each vertex.

Observations

- 1 If Γ is the 3-graph of cyclic triangles in any tournament, then the links Γ_v of its vertices are bipartite.
- 2 If Γ is *any* 3-graph with bipartite links, then Γ contains no $K_4^{(3)} - e$.

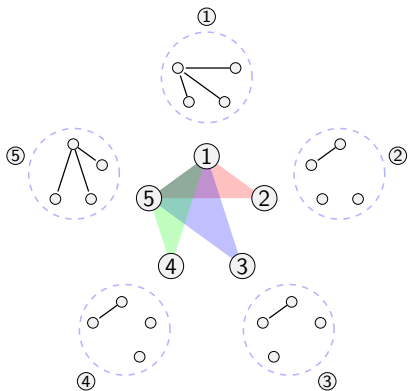


Figure: The links of each vertex.

Observations

- 1 If Γ is the 3-graph of cyclic triangles in any tournament, then the links Γ_v of its vertices are bipartite.
- 2 If Γ is *any* 3-graph with bipartite links, then Γ contains no $K_4^{(3)} - e$.
- 3 We can reproduce the Erdős-Hajnal lower bound by taking Γ to be a 3-graph with random bipartite links.

Modified Construction

Let Γ be a random 3-graph on N vertices specified by N bipartitions $U_v \cup W_v = V(G) \setminus \{v\}$ indexed by the vertices v . A triple $\{u, v, w\}$ is an edge of Γ iff v and w are on the opposite sides of the bipartition for u , u and w are on opposite sides of the bipartition for v , and u and v are on opposite sides of the bipartition for w .

Modified Construction

Let Γ be a random 3-graph on N vertices specified by N bipartitions $U_v \cup W_v = V(G) \setminus \{v\}$ indexed by the vertices v . A triple $\{u, v, w\}$ is an edge of Γ iff v and w are on the opposite sides of the bipartition for u , u and w are on opposite sides of the bipartition for v , and u and v are on opposite sides of the bipartition for w .

Lemma

If the bipartitions are chosen uniformly at random, then w.h.p. Γ is $(K_4^{(3)} - e)$ -free, has edge density $1/8 + o(1)$, and independence number $\Theta(\log N)$.

Random bipartite links

Modified Construction

Let Γ be a random 3-graph on N vertices specified by N bipartitions $U_v \cup W_v = V(G) \setminus \{v\}$ indexed by the vertices v . A triple $\{u, v, w\}$ is an edge of Γ iff v and w are on the opposite sides of the bipartition for u , u and w are on opposite sides of the bipartition for v , and u and v are on opposite sides of the bipartition for w .

Lemma

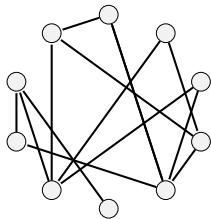
If the bipartitions are chosen uniformly at random, then w.h.p. Γ is $(K_4^{(3)} - e)$ -free, has edge density $1/8 + o(1)$, and independence number $\Theta(\log N)$.

But we want independence number $n = O(\log N / \log \log N)$ to get $N = 2^{\Omega(n \log n)}$, so the links can't be bipartite.

Our Construction

Naive random construction

①



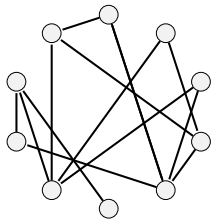
Problem

Needs to be very sparse.

Our Construction

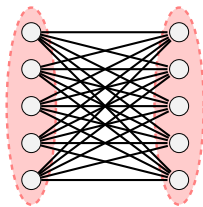
Naive random construction

①



Random bipartite links

①



Problem

Needs to be very sparse.

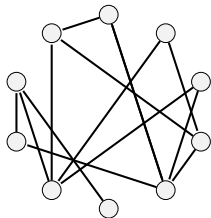
Problem

Large independence number due to bipartition.

Our Construction

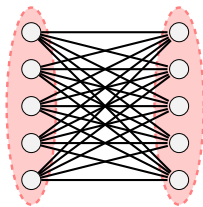
Naive random construction

①



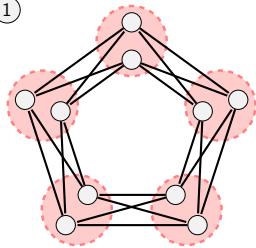
Random bipartite links

①



Random triangle-free links

①



Problem

Needs to be very sparse.

Problem

Large independence number due to bipartition.

Our construction

Links are random blowups of a small triangle-free graph.

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

Construction

Fix an auxiliary graph A on $m = n^C$ vertices, which is triangle-free and has edge density $p = m^{-2/3}$.

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

Construction

Fix an auxiliary graph A on $m = n^C$ vertices, which is triangle-free and has edge density $p = m^{-2/3}$.

Let Γ be the 3-graph on $N = 2^{cn \log n}$ vertices specified by a map $\chi : V(\Gamma)^2 \rightarrow V(A)$. A triple $\{u, v, w\}$ is an edge of Γ iff $\chi(u, v) \sim \chi(u, w)$, $\chi(v, u) \sim \chi(v, w)$, and $\chi(w, u) \sim \chi(w, v)$.

Proof of main result

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

Construction

Fix an auxiliary graph A on $m = n^C$ vertices, which is triangle-free and has edge density $p = m^{-2/3}$.

Let Γ be the 3-graph on $N = 2^{cn \log n}$ vertices specified by a map $\chi : V(\Gamma)^2 \rightarrow V(A)$. A triple $\{u, v, w\}$ is an edge of Γ iff $\chi(u, v) \sim \chi(u, w)$, $\chi(v, u) \sim \chi(v, w)$, and $\chi(w, u) \sim \chi(w, v)$.

Choose χ uniformly at random. For each v , the link Γ_v will be a subgraph of a blowup of A , so it is triangle-free.

Proof of main result

Theorem (Fox-H. 2019)

$$r(K_4^{(3)} - e, K_n^{(3)}) = 2^{\Theta(n \log n)}.$$

Construction

Fix an auxiliary graph A on $m = n^C$ vertices, which is triangle-free and has edge density $p = m^{-2/3}$.

Let Γ be the 3-graph on $N = 2^{cn \log n}$ vertices specified by a map $\chi : V(\Gamma)^2 \rightarrow V(A)$. A triple $\{u, v, w\}$ is an edge of Γ iff $\chi(u, v) \sim \chi(u, w)$, $\chi(v, u) \sim \chi(v, w)$, and $\chi(w, u) \sim \chi(w, v)$.

Choose χ uniformly at random. For each v , the link Γ_v will be a subgraph of a blowup of A , so it is triangle-free.

Lemma (Hard part)

If Γ is the random 3-graph described above, the independence number of Γ is less than $n = O(\log N / \log \log N)$.

Definition

If G is a k -graph, the *link hypergraph* L_G of G is the $(k + 1)$ -graph on $V(G) \cup \{v\}$ (for a new vertex v) whose edges come from inserting v into all the edges of G .

General link hypergraphs

Definition

If G is a k -graph, the *link hypergraph* L_G of G is the $(k + 1)$ -graph on $V(G) \cup \{v\}$ (for a new vertex v) whose edges come from inserting v into all the edges of G .

Theorem (Conlon-Fox-Sudakov 2010)

If G is bipartite, then $r(L_G, K_n^{(3)}) = n^{\Theta(1)}$.

If G is non-bipartite, then $r(L_G, K_n^{(3)}) = 2^{\Omega(n)}$.

General link hypergraphs

Definition

If G is a k -graph, the *link hypergraph* L_G of G is the $(k + 1)$ -graph on $V(G) \cup \{v\}$ (for a new vertex v) whose edges come from inserting v into all the edges of G .

Theorem (Conlon-Fox-Sudakov 2010)

If G is bipartite, then $r(L_G, K_n^{(3)}) = n^{\Theta(1)}$.

If G is non-bipartite, then $r(L_G, K_n^{(3)}) = 2^{\Omega(n)}$.

Theorem (Fox-H. 2019)

If G is non-bipartite, then $r(L_G, K_n^{(3)}) = 2^{\Theta(n \log n)}$.

General link hypergraphs

Definition

If G is a k -graph, the *link hypergraph* L_G of G is the $(k+1)$ -graph on $V(G) \cup \{v\}$ (for a new vertex v) whose edges come from inserting v into all the edges of G .

Theorem (Conlon-Fox-Sudakov 2010)

If G is bipartite, then $r(L_G, K_n^{(3)}) = n^{\Theta(1)}$.

If G is non-bipartite, then $r(L_G, K_n^{(3)}) = 2^{\Omega(n)}$.

Theorem (Fox-H. 2019)

If G is non-bipartite, then $r(L_G, K_n^{(3)}) = 2^{\Theta(n \log n)}$.

Question

How do the implicit constants depend on G ?

Theorem (Fox-H. 2019)

For all $s \geq 3$ and $n \geq 1$, we have

$$r(L_{K_s}, K_{n,n,n}^{(3)}) = \binom{n+s}{s}^{\Theta(n)}.$$

Theorem (Fox-H. 2019)

For all $s \geq 3$ and $n \geq 1$, we have

$$r(L_{K_s}, K_{n,n,n}^{(3)}) = \binom{n+s}{s}^{\Theta(n)}.$$

Corollary

If $n \geq 3$, then

$$r(L_{K_n}, K_{n,n,n}^{(3)}) = 2^{\Theta(n^2)}.$$

Diagonal Ramsey numbers via link hypergraphs

Theorem (Fox-H. 2019)

For all $s \geq 3$ and $n \geq 1$, we have

$$r(L_{K_s}, K_{n,n,n}^{(3)}) = \binom{n+s}{s}^{\Theta(n)}.$$

Corollary

If $n \geq 3$, then

$$r(L_{K_n}, K_{n,n,n}^{(3)}) = 2^{\Theta(n^2)}.$$

Remark

This gives a very “sparse” proof that $r(K_n^{(3)}, K_n^{(3)}) = 2^{\Omega(n^2)}$, and suggests that the diagonal Ramsey number should be much bigger.

Ramsey transition thresholds

Definition

Let $r_k(s, t; n)$ be the minimum N such that in any k -graph on N vertices, either some s vertices span t edges, or else there is an independent set of size n .

Ramsey transition thresholds

Definition

Let $r_k(s, t; n)$ be the minimum N such that in any k -graph on N vertices, either some s vertices span t edges, or else there is an independent set of size n .

Examples

Thus $r_3(4, 4; n) = r(K_4^{(3)}, K_n^{(3)})$ and $r_3(4, 3; n) = r(K_4^{(3)} - e, K_n^{(3)})$.

Ramsey transition thresholds

Definition

Let $r_k(s, t; n)$ be the minimum N such that in any k -graph on N vertices, either some s vertices span t edges, or else there is an independent set of size n .

Examples

Thus $r_3(4, 4; n) = r(K_4^{(3)}, K_n^{(3)})$ and $r_3(4, 3; n) = r(K_4^{(3)} - e, K_n^{(3)})$.

Conjecture (Erdős-Hajnal 1972)

For every $s > k \geq 3$, there exists a unique $t = h_1^{(k)}(s)$ such that $r_k(s, t - 1; n)$ is polynomial in n and $r_k(s, t; n)$ is exponential in n .

Ramsey transition thresholds

Definition

Let $r_k(s, t; n)$ be the minimum N such that in any k -graph on N vertices, either some s vertices span t edges, or else there is an independent set of size n .

Examples

Thus $r_3(4, 4; n) = r(K_4^{(3)}, K_n^{(3)})$ and $r_3(4, 3; n) = r(K_4^{(3)} - e, K_n^{(3)})$.

Conjecture (Erdős-Hajnal 1972)

For every $s > k \geq 3$, there exists a unique $t = h_1^{(k)}(s)$ such that $r_k(s, t - 1; n)$ is polynomial in n and $r_k(s, t; n)$ is exponential in n . In general, for every $1 \leq i \leq k - 2$ there is a unique $t = h_i^{(k)}(s)$ such that $r_k(s, t - 1; n)$ has tower height i and $r_k(s, t; n)$ has tower height $i + 1$ (as a function of n).

The polynomial-to-exponential transition

Theorem (Conlon-Fox-Sudakov 2010)

For infinitely many s , $h_1^{(3)}(s)$ exists and $h_1^{(3)}(s) - 1 = T(s)$ which is the maximum number of cyclic triangles in a tournament on s vertices.

The polynomial-to-exponential transition

Theorem (Conlon-Fox-Sudakov 2010)

For infinitely many s , $h_1^{(3)}(s)$ exists and $h_1^{(3)}(s) - 1 = T(s)$ which is the maximum number of cyclic triangles in a tournament on s vertices.

Theorem (Mubayi-Razborov 2019)

For all $s > k \geq 4$, $h_1^{(k)}(s)$ exists and $h_1^{(k)}(s) - 1 = g^{(k)}(s)$ which is the maximum number of ordered rainbow tournaments on k vertices in an ordered $\binom{k}{2}$ -edge-colored tournament on s vertices.

The polynomial-to-exponential transition

Theorem (Conlon-Fox-Sudakov 2010)

For infinitely many s , $h_1^{(3)}(s)$ exists and $h_1^{(3)}(s) - 1 = T(s)$ which is the maximum number of cyclic triangles in a tournament on s vertices.

Theorem (Mubayi-Razborov 2019)

For all $s > k \geq 4$, $h_1^{(k)}(s)$ exists and $h_1^{(k)}(s) - 1 = g^{(k)}(s)$ which is the maximum number of ordered rainbow tournaments on k vertices in an ordered $\binom{k}{2}$ -edge-colored tournament on s vertices.

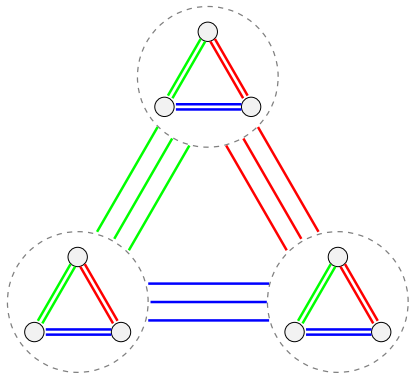
Both results relate thresholds to the general problem of *inducibility*: given a (possibly colored and/or directed) graph on k vertices, what is the maximum number of induced copies of it in a graph on s vertices?

Open problem: inducibility for $k = 3$

Conjecture (Erdős-Hajnal 1972)

For all $s \geq 3$, the maximum number of ordered rainbow triangles in a 3-edge-coloring of K_s is $g^{(3)}(s)$, where $g^{(3)}(s) = 0$ if $s < 3$ and otherwise

$$g^{(3)}(s) = \max_{a+b+c=s} \{g^{(3)}(a) + g^{(3)}(b) + g^{(3)}(c) + abc\}.$$



Theorem (Fox-H. 2019)

We have $r_3(4, 2; n) = n^{\Theta(1)}$ but $r_3(4, 3; n) = 2^{\Omega(n \log n)}$.

Skipping the exponential order

Theorem (Fox-H. 2019)

We have $r_3(4, 2; n) = n^{\Theta(1)}$ but $r_3(4, 3; n) = 2^{\Omega(n \log n)}$.

Theorem (Fox-H. 2019)

For s large and $.26 \binom{s}{3} \leq t \leq .46 \binom{s}{3}$, $r_3(s, t; n) = 2^{\Theta(n \log n)}$.

Skipping the exponential order

Theorem (Fox-H. 2019)

We have $r_3(4, 2; n) = n^{\Theta(1)}$ but $r_3(4, 3; n) = 2^{\Omega(n \log n)}$.

Theorem (Fox-H. 2019)

For s large and $.26 \binom{s}{3} \leq t \leq .46 \binom{s}{3}$, $r_3(s, t; n) = 2^{\Theta(n \log n)}$.

Conjecture

For all $s \geq 4$, there exists t for which $r_3(s, t - 1; n) = n^{\Theta(1)}$ and $r_3(s, t; n) = 2^{\Omega(n \log n)}$.

Thank you!

