

Selected Nonlinear Optimization Applications

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Chapters 7.2, 11.4, and Wikipedia

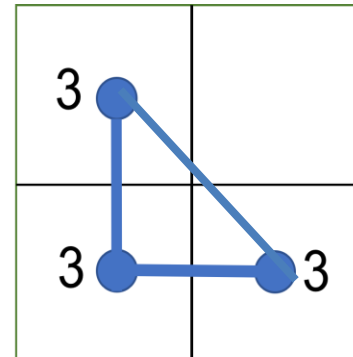
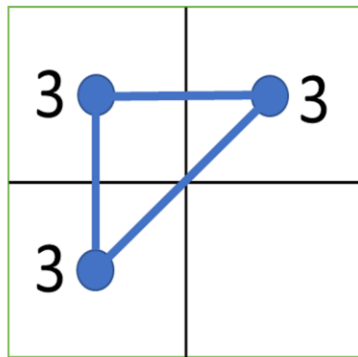
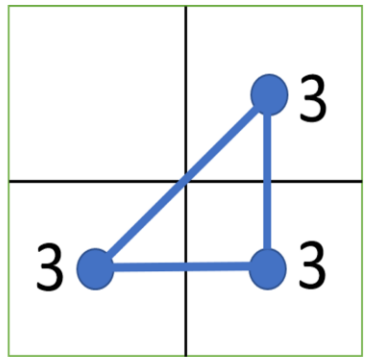
The Linear Regression: Least Squares Model

For the Business-or-Personal problem, we now minimize the sum of the squared errors (between predicted personal remittances and actual personal remittances).

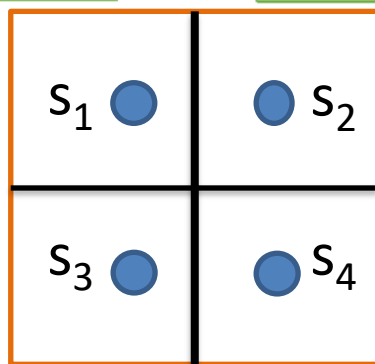
$$\begin{aligned} \text{Min} \quad & \sum_i \left(\sum_j a_{ij} x_j - b_i \right)^2 \\ \text{s.t.} \quad & 0 \leq x_j \leq 1, \forall j. \end{aligned}$$

- Let x_j be such a probability that a transaction is personal for industry code j
- $a_{i,j}$ – transaction amount for account i and industry code j
- b_i – amount paid by personal remit for account i
- $\sum_j a_{i,j} x_j$ – the expected personal expenses for account i
- We'd like to choose x_j such that $\sum_j a_{i,j} x_j$ matches b_i for ALL i
- This model is called **Quadratic Optimization**
- **Convex? 1st Order Optimality Conditions? Sufficient?**

The Wasserstein Barycenter Optimization



Three possible demand distribution scenario of 4 cities



Constraints:

$$s_1 + s_2 + s_3 + s_4 = 9$$

$$(s_1, s_2, s_3, s_4) \geq 0$$

$$C = \begin{pmatrix} 0 & 1 & 1 & 2 \\ 1 & 0 & 2 & 1 \\ 1 & 2 & 0 & 1 \\ 2 & 1 & 1 & 0 \end{pmatrix}$$

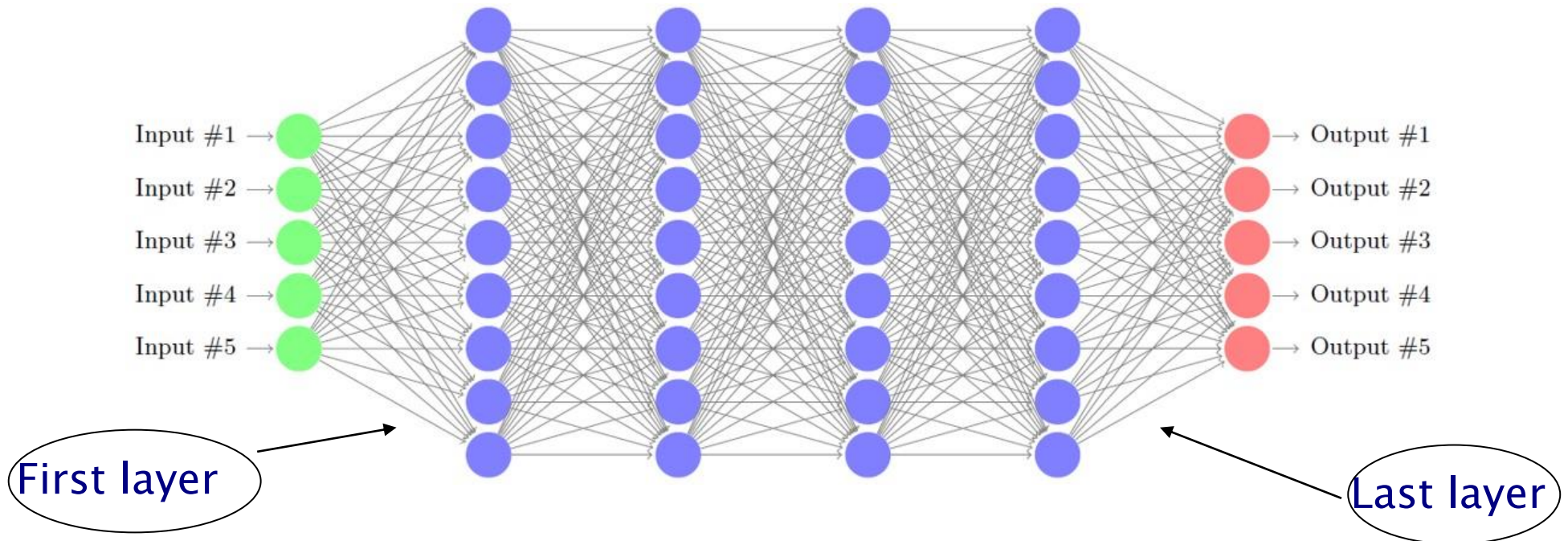
$$\text{minimize } WD_l(\mathbf{s}) + WD_r(\mathbf{s}) + WD_r(\mathbf{s})$$

$$\text{s.t. } s_1 + s_2 + s_3 + s_4 = 9, \mathbf{s} \geq \mathbf{0}$$

Wasserstein-Distance function, $WD(\mathbf{s})$, is an **implicit nonlinear but convex** function defined by the minimum value of a transportation minimization problem from supply inventory distribution \mathbf{s} .

This is a linearly constrained **convex nonlinear** optimization problem.

Nonlinear Optimization in Deep Learning



$$\text{minimize } L(\text{ReLU}(a_i(w_{i,j})))$$

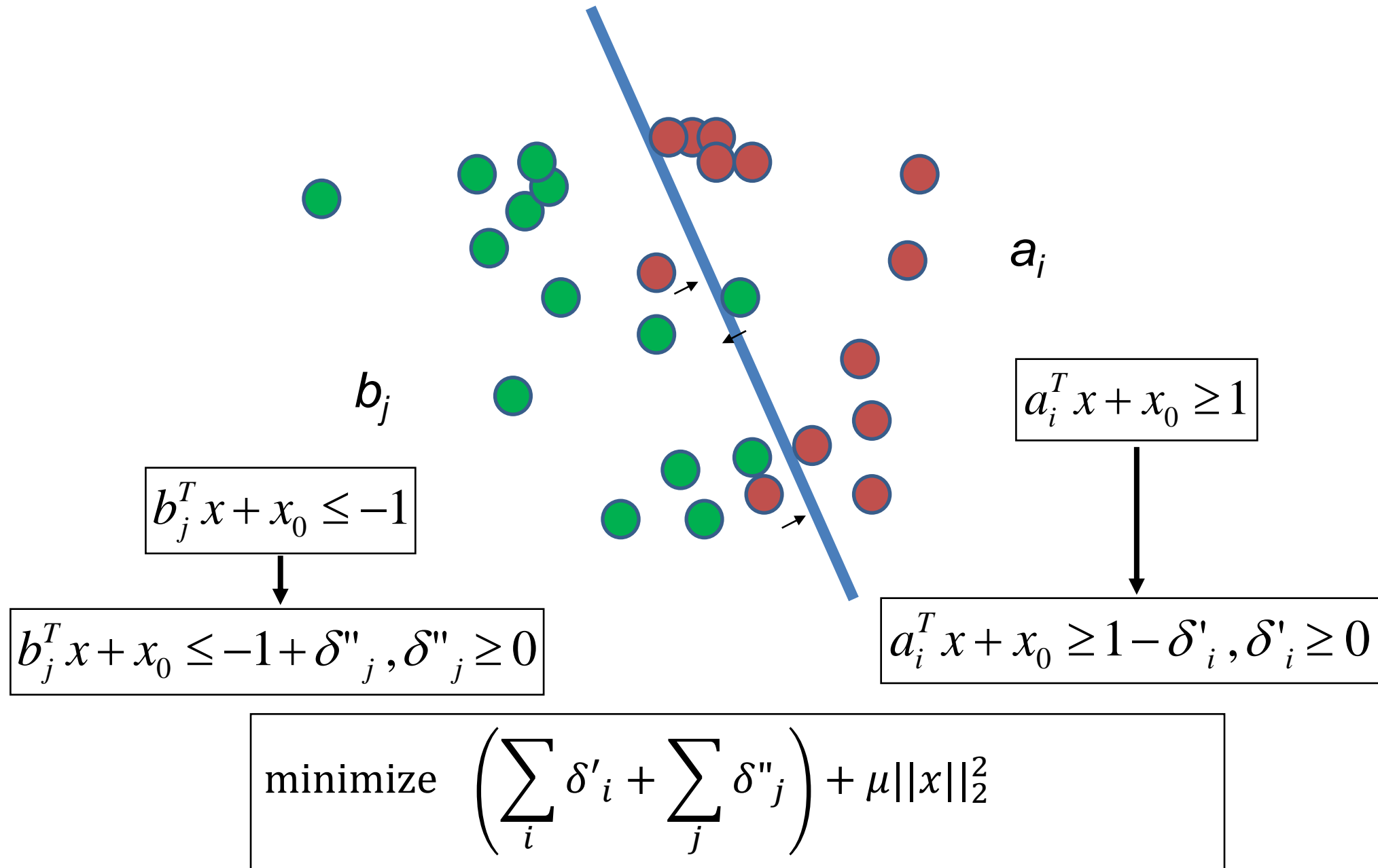
where $w_{i,j}$ is the weight variable at layer i and edge j ,

$(a_i(w_{i,j}))$ is a linear function of weights

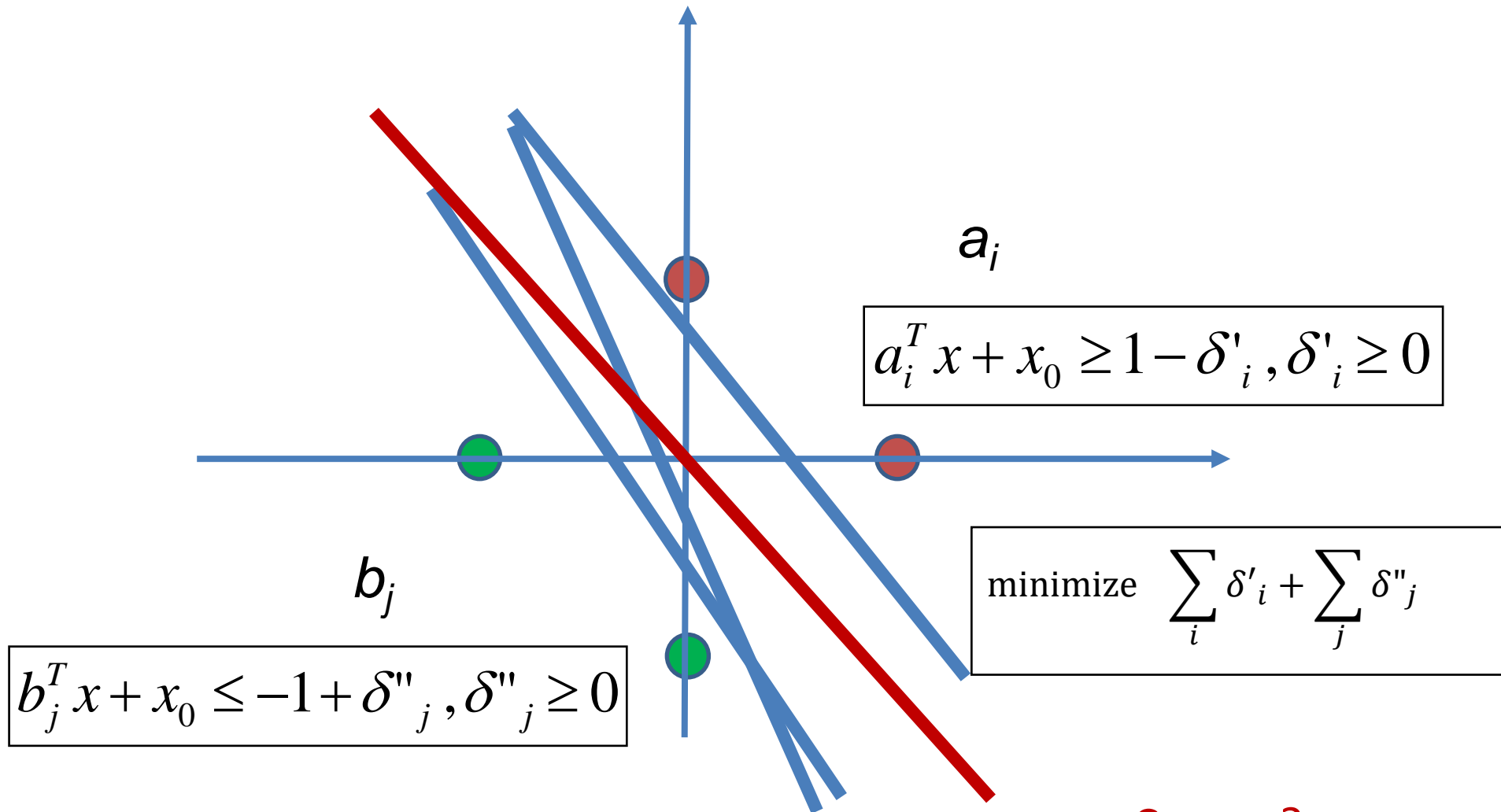
$\text{ReLU}(\cdot)$ is the function $\max\{\cdot, 0\}$, called **Rectified Linear Unit** function (a type of neural behavior) in Deep Learning. The problem is typically

Nonlinear and Nonconvex.

SVM with Quadratic Optimization I



SVM with Quadratic Optimization II



$$\text{minimize } \left(\sum_i \delta'_i + \sum_j \delta''_j \right) + 0.1 \|x\|_2^2$$

Convex?
 1st Order
 Optimality
 Conditions?
 Sufficient?

SVM vs Logistic Regression: Likelihood Probability

Given message \mathbf{a}_i , according to the **logistic model**, the probability that it's a SPAM is represented by

$$\frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)}$$

Thus, for the training data, we like to determine x_0 and \mathbf{x} from a set of **training data (some spam some not)** such that


$$\frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \sim \begin{cases} 1, & \text{spam} \\ 0, & \text{not.} \end{cases}$$

The probability to give a "right answer" for all training messages is

$$\left[\prod_{i:\text{spam}} \frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \right] \left[\prod_{i:\text{not}} \left[1 - \frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \right] \right]$$

Logistic Regression: Max-Likelihood Probability

Thus, we like to determine x_0 and \mathbf{x} to maximize

$$\left[\prod_{i:spam} \frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \right] \left[\prod_{i:not} \left[1 - \frac{\exp(\mathbf{a}_i^T \mathbf{x} + x_0)}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \right] \right]$$


$$\left[\prod_{i:spam} \frac{1}{1 + \exp(-\mathbf{a}_i^T \mathbf{x} - x_0)} \right] \left[\prod_{i:not} \frac{1}{1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)} \right]$$

which is equivalent to maximize the **log-likelihood** probability

$$-\sum_{i:spam} \log(1 + \exp(-\mathbf{a}_i^T \mathbf{x} - x_0)) - \sum_{i:not} \log(1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)).$$

Or to minimize the convex (!) **log-logistic-loss** (with possible L_2 regularization term $\|\mathbf{x}\|^2$ added into the objective)

$$\sum_{i:spam} \log(1 + \exp(-\mathbf{a}_i^T \mathbf{x} - x_0)) + \sum_{i:not} \log(1 + \exp(\mathbf{a}_i^T \mathbf{x} + x_0)).$$

Convex?
1st Order
Optimality
Conditions?
Sufficient?

Convexity of log-exponential-sum

$$f(x_1, x_2) = \log(e^{x_1} + e^{x_2})$$

$$\nabla f(x_1, x_2) = \begin{pmatrix} \frac{e^{x_1}}{e^{x_1} + e^{x_2}} \\ \frac{e^{x_2}}{e^{x_1} + e^{x_2}} \end{pmatrix}$$

$$\nabla^2 f(x_1, x_2) = \frac{1}{(e^{x_1} + e^{x_2})^2} \begin{pmatrix} e^{x_1+x_2} & -e^{x_1+x_2} \\ -e^{x_1+x_2} & e^{x_1+x_2} \end{pmatrix}$$

The Hessian matrix is PSD everywhere so that LR is an
Unconstrained Convex Program

Geometric Optimization and Convexification

$$\begin{aligned} \min \quad & x \cdot y + y \cdot z + z \cdot x \\ \text{s.t.} \quad & x \cdot y \cdot z = 1, \\ & (x, y, z) \geq 0 \end{aligned}$$



$$\begin{aligned} \min \quad & e^{u+v} + e^{v+w} + e^{w+u} \\ \text{s.t.} \quad & e^{u+v+w} = 1 \text{ or } u + v + w = 0. \end{aligned}$$

Let $u = \ln(x)$, $v = \ln(y)$, $w = \ln(z)$, then

This a **Linearly Constrained (Convex) Optimization Program**

Dynamic Optimal Pricing

p : price decision

$q(p) = \exp(a \log(p) + b)$: demand volume function

a : elasticity coefficient (< 0)

b : fixed and/or extenality coefficient

c : unit cost

Profit function: $(p-c) \cdot q(p) = (p-c) \cdot \exp(a \log(p) + b)$

Optimal price:

$$p^* = \begin{cases} \frac{ac}{a+1} & \text{if } a < -1 \\ c & \text{otherwise} \end{cases} : \text{optimal price to maximiza profit}$$

Note that the optimal price does not depend on b .



Use historical data and regression method to learn the **demand function coefficients**, and then compute the optimal price and fixed it for the remaining time periods.

Dynamic Pricing: update demand coefficients using new data and recalculate optimal price.

Online Combinatorial Auction I

Order fill	Price Limit π	Quantity Limit q	Argentina	Brazil	Italy	Germany	France
x1	0.75	10	1	1	1		
x2	0.35	5				1	
x3	0.40	10	1		1		1
x4	0.95	10	1	1	1	1	
x5	0.75	5		1		1	

$$\max_{\mathbf{x}, w} \sum_{j=1}^n \pi_j x_j - w$$

$$\text{s.t.} \quad \sum_{j=1}^n \mathbf{a}_j x_j - \mathbf{1}w \leq \mathbf{0},$$

$$0 \leq x_j \leq q_j, \forall j$$

π_j : the j th bidding price

\mathbf{a}_j : the j th bidding vector;

q_j : the j th bidding share up limit;

**Now Consider the online decisions
And the nonlinear optimization model**

$$\max_{\{x, w, \mathbf{s}\}} \sum_{j=1}^n \pi_j x_j - w + \sum_{i=1}^m u_i(s_i)$$

$$\text{s.t.} \quad \sum_{j=1}^n \mathbf{a}_j x_j - \mathbf{1}w + \mathbf{s} = \mathbf{0},$$

$$0 \leq x_j \leq q_j, \mathbf{s} \geq \mathbf{0}.$$

$u(s)$: **Nonlinear Concave Value Function**

Logarithmic function with a positive weight μ

$$u(s) = \mu \ln(s)$$

Exponential function with a positive weight μ

$$u(s) = \mu(1 - \exp(-s/\mu))$$

Sequential Convex Programming Mechanism: OCA II

$$\begin{aligned}
 \max \quad & \sum_{j=1}^n \pi_j x_j - w \\
 \text{s.t.} \quad & \sum_{j=1}^n \mathbf{a}_j x_j - \mathbf{1}w \leq \mathbf{0}, \\
 & 0 \leq x_j \leq q_j, \forall j
 \end{aligned}$$

Suppose the order fill decision x_j^* from $j=1$ to k have been made. Now the k th bidder arrives and we solve the following convex optimization problem

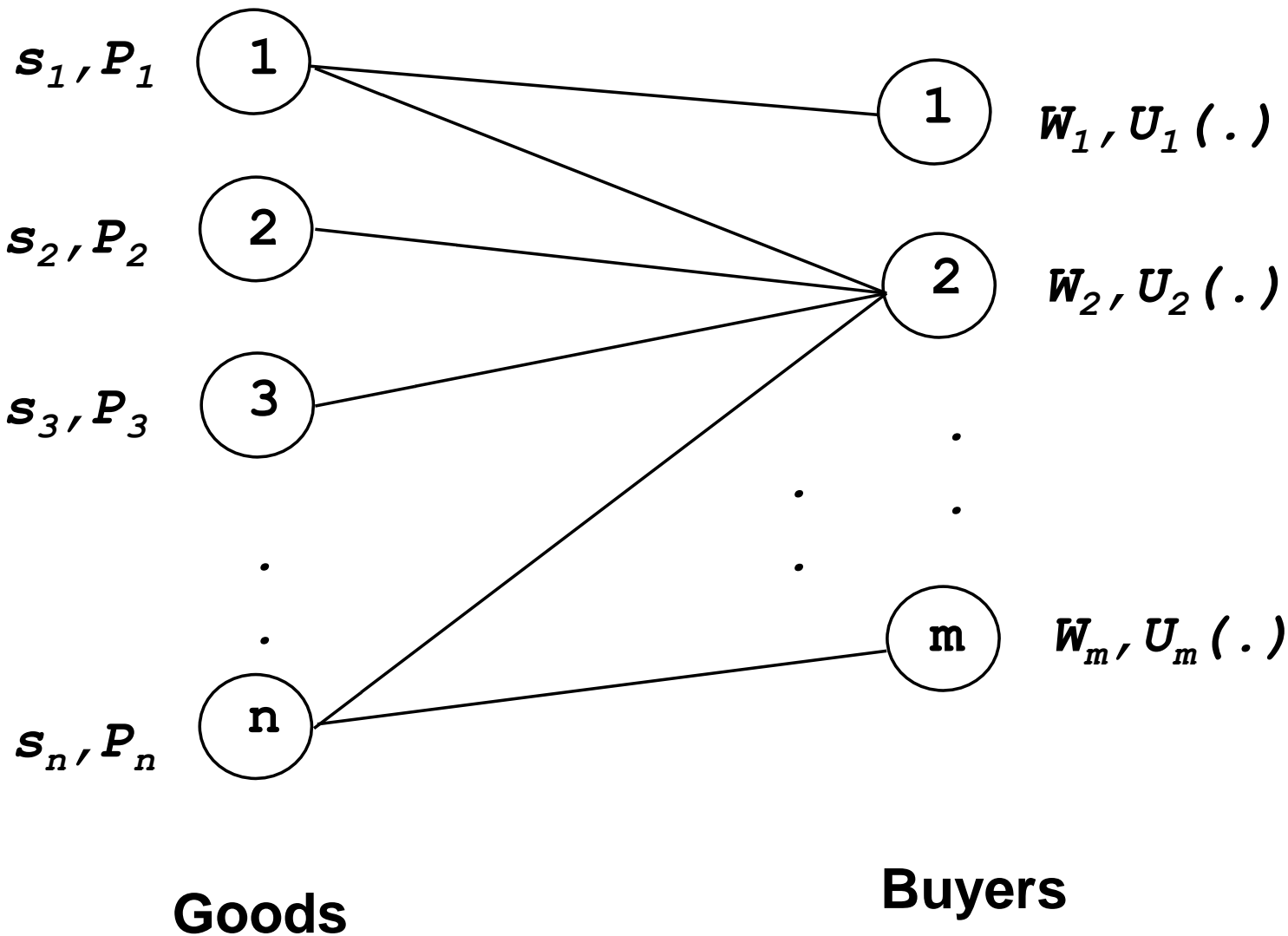
$$\begin{aligned}
 \max_{\{x, w, \mathbf{s}\}} \quad & \pi_k x - w + \sum_{i=1}^m u_i(s_i) \\
 \text{s.t.} \quad & \mathbf{a}_k x - \mathbf{1}w + \mathbf{s} = - \sum_{j=1}^{k-1} \mathbf{a}_j x_j^*, \\
 & 0 \leq x \leq q_k, \quad \mathbf{s} \geq \mathbf{0}.
 \end{aligned}$$

Slack variables, or shares held by us.

Value functions to hold own shares

Outstanding shares already sold to bidders.

Fisher's Equilibrium Market Model



- Buyers have a money budget (w_i) to buy goods and maximize their individual utility functions $u(\cdot)$;
- Producers have goods (s_j) to sell for money

• The equilibrium price is an assignment of prices (p_j) to goods so that every buyer could buy a maximal bundle of goods and also clear the market, meaning that all the money are spent and all goods are sold.

Each Buyer's Maximization and Equilibrium Price

Buyer $i = 1, \dots, m$ optimization problem for given prices p_j , $j = 1, \dots, n$.

$$\begin{aligned} \max_{\mathbf{x}} \quad & \mathbf{u}_i^T \mathbf{x}_i := \sum_j u_{ij} x_{ij} \\ \text{s.t.} \quad & \mathbf{p}^T \mathbf{x}_i := \sum_j p_j x_{ij} \leq w_i, \\ & x_{ij} \geq 0, \quad \forall \quad j, \end{aligned}$$

The **equilibrium price vector** \mathbf{p} is the one to make

$$\sum_i x_{ij}^* = s_j$$

where \mathbf{x}_i^* is an optimal solution vector of the above problem with given prices \mathbf{p} for all $i = 1, \dots, m$.

Example of Fisher's Model and Equilibrium Conditions

Buyer 1 and 2's optimization problems for given prices p_x, p_y on two goods x and y, each has one unit on the market.

$$\begin{aligned} \max \quad & 2x_1 + y_1 \\ \text{s.t.} \quad & p_x \cdot x_1 + p_y \cdot y_1 \leq 5, \\ & x_1, y_1 \geq 0; \end{aligned}$$

$$\begin{aligned} \max \quad & 3x_2 + y_2 \\ \text{s.t.} \quad & p_x \cdot x_2 + p_y \cdot y_2 \leq 8, \\ & x_2, y_2 \geq 0. \end{aligned}$$

$$\begin{aligned} p_x \lambda_1 - 2 &\geq 0 \quad \wedge \quad x_1 \\ p_y \lambda_1 - 1 &\geq 0 \quad \wedge \quad y_1 & x_1 + x_2 = 1 \\ & & y_1 + y_2 = 1 \\ p_x \lambda_2 - 3 &\geq 0 \quad \wedge \quad x_2 \\ p_y \lambda_2 - 1 &\geq 0 \quad \wedge \quad y_2 \end{aligned}$$

$$\begin{aligned} p_x &= \frac{26}{3}, \quad p_y = \frac{13}{3} \\ x_1 &= \frac{6}{78}, \quad y_1 = 1, \quad x_2 = \frac{72}{78}, \quad y_2 = 0 \end{aligned}$$

satisfy all these conditions so that the prices are equilibrium prices

Aggregate Social Optimization Problem

$$\begin{aligned} \max \quad & \sum_i w_i \log(\mathbf{u}_i^T \mathbf{x}_i) \\ \text{s.t.} \quad & \sum_i x_{ij} = s_j, \quad \forall j=1, \dots, n \\ & x_{ij} \geq 0, \quad \forall i, j, \end{aligned}$$

Theorem (*Eisenberg and Gale 1959*) *The optimal Lagrange multiplier vector of the equality constraints of the social optimization problem is the equilibrium price vector.*

Aggregate Social Problem for the Example

$$\begin{aligned} \max \quad & 5 \log(2x_1 + y_1) + 8 \log(3x_2 + y_2) \\ \text{s.t.} \quad & x_1 + x_2 = 1, \\ & y_1 + y_2 = 1, \\ & (x_1, x_2, y_1, y_2) \geq 0. \end{aligned}$$

$$\begin{aligned} \frac{10}{2x_1 + y_1} &\leq p_x \wedge x_1 & \frac{5}{2x_1 + y_1} &\leq p_y \wedge y_1 \\ \frac{24}{3x_2 + y_2} &\leq p_x \wedge x_2 & \frac{8}{3x_2 + y_2} &\leq p_y \wedge y_2 \end{aligned}$$

The optimality conditions of the social optimization problem.

By assign

$$5\lambda_1 = 2x_1 + y_1$$

$$8\lambda_2 = 3x_2 + y_2$$

these conditions coincides The equilibrium conditions on Slide 16.

Sensor Network Localization I

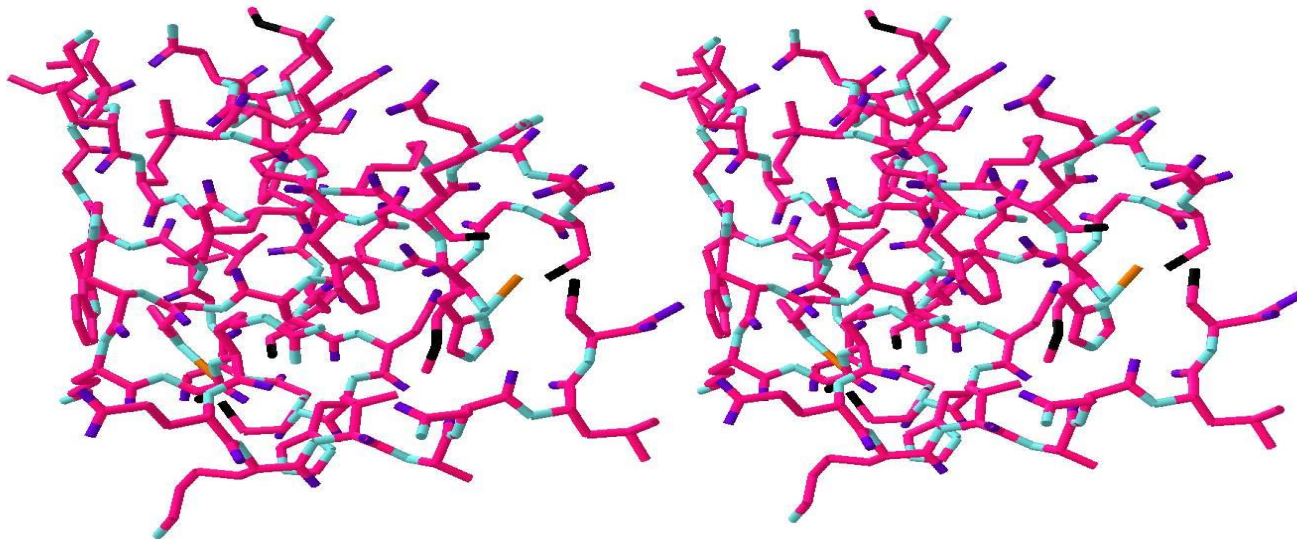
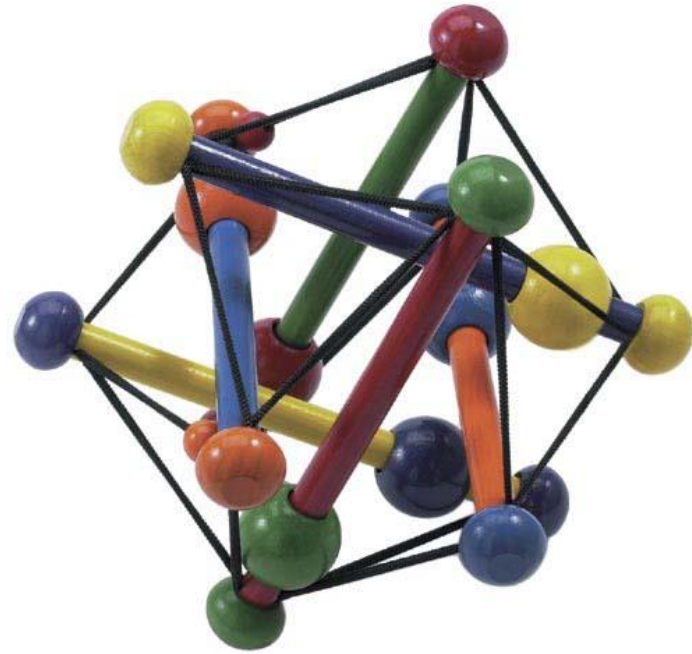
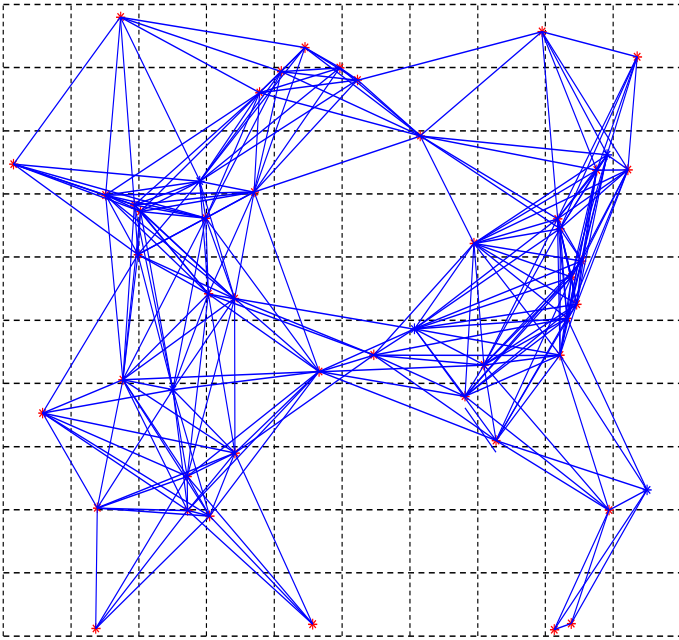
Given a graph $G = (V, E)$ and sets of partial **distance measurements**, say $\{d_{ij} : (i, j) \in E\}$, the goal is to compute a **realization** of G in the **Euclidean space \mathbf{R}^d** for a **given low dimension d** , i.e.

- to place the nodes/vertices of G in \mathbf{R}^d such that
- the **Euclidean distance** between every pair of adjacent vertices $(i, j) \in E$ equals the measurements $d_{ij} \in E$.

In general the localization may not be fixed since the configuration can rotate and translate. Thus, we assume that the positions of **$(d+1)$** sensors are known, and they called anchors.

This problem has wide applications ...

Sensor Network Localization II



Sensor Network Localization III

Find d -dimensional points/vectors $\mathbf{x}_j, j=d+2, d+3, \dots, n$, such that

$$\|x_i - x_j\| = d_{ij} \quad \forall (i, j) \in E$$

Anchors : $x_i = a_i, \quad i = 1, 2, \dots, d + 1$

This is a system of **quadratic equations** (after square both sides) and nonconvex, in contrast to a system of linear equations.

Does the system have a solution/localization of all \mathbf{x}_j 's? Is the solution/localization **unique**? Is there a **certification** for a solution to make it **reliable or trustworthy**? Is the system **partially** localizable with certification?

To get something tractable, we can consider optimization formulation or convex relaxation approaches.

SNL: Nonlinear Least Squares and Convex Relaxation

One can form a nonlinear least square minimization problem

$$\text{minimize } \sum_{(ij) \in E} (\|x_i - x_j\|^2 - (d_{ij})^2)^2$$

$$\text{s.t. } x_i = a_i, \text{ for } i=1, \dots, d+1$$

This remains nonconvex **quartic polynomial minimization** (with many local minimums)

We would use various algorithms to tackle the problem, including some convex relaxation approaches, which would be discussed later in the class.

$$\text{minimize } 0^T x:$$

$$\text{s.t. } \|x_i - x_j\| \leq (d_{ij}) \text{ for } (i,j) \in E$$

$$\text{where } x_i = a_i, \text{ for } i=1, \dots, d+1$$

Convex?
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