Introduction to Optimization

Yinyu Ye Department of Management Science and Engineering Stanford University Stanford, CA 94305, U.S.A.

https://canvas.stanford.edu/courses/142617

Read Chapter 1.1, 1.2, 2.1, 2.2, Appendices A, B&D in Text-Book (hard copies would be available in the Book Store)

1st Day Questions

- My CA team: Aldo, Luc, Chunlin, Mingxi
- Websites: Canvas

https://canvas.stanford.edu/courses/142617

- 4 Homework assignments, 1 Take-Home Midterm, 1 team project
	- $-40\% * H + 30\% * M + 30\% * P$
	- No difference on taking 3 or 4 units
- No formula for cutoff between A/B etc.
- The more fun we all have, the more A's we will give out.
- Textbook: Linear and Nonlinear Programming (LY 4-5th edition, posted in Canvas)
- This is 111X/211X; if never had calculus and linear algebra classes, take 111/211
- The software use will help: Solvers in Matlab, R, Python or other public free software. It is mostly a "PAPER AND PENCIL" class!
- Form a "diversified" study group
- Selected Friday's problem sessions (will be taped)
- Students with OAE, extra two day for the exam

Mathematical Optimization Model

- Often consider the common quantitative model of data/decision/management science & engineering:
	- Maximize or Minimize *f(x) for* all $x \in$ some set X
- Decision variables *x*, Objective function *f(x)*, Constraint set X
- Applications in:
	- Applied Science, Engineering, Economics, Finance, Medicine, Statistics, Business
	- General Decision and Policy Making
- The famous Eighteenth Century Swiss mathematician and physicist Leonhard Euler (1707-1783) proclaimed that "...nothing at all takes place in the Universe in which some rule of maximum or minimum does not appear."

The Prototypical Optimization Problem

The Function could be:

 $x_1 + 2x_2, x^2 + 2xy + 2y^2, xln(x) + e^y, \sqrt{x} + max\{x, y\}$, etc

Linear Programming/Optimization: all functions are linear/affine

Yinyu Ye, Stanford, MS&E211 Lecture Notes #1 4 4

Linear Programming

- Why do we study LP's
	- Not just because solving non-linear problems are difficult
	- But because real-world problems are often formulated as linear equations and inequalities
		- Either because they indeed are linear
		- Or because it is unclear how to represent them and linear is an intuitive compromise
	- A stepping stone for solving more complicated nonlinear optimization problems, which you would see later.

… and National Medal of Science

3 Main Categories in Optimization Covered in this Course

- **Linear Optimization** (Programming)
	- Search Algorithms
	- Simplex and Interior-Point Algorithms
- **Unconstrained Nonlinear Optimization**
	- 1 st order methods, gradient method
	- 2nd order methods, Newton
- **Constrained Nonlinear Optimization**
	- 1 st order, gradient projection, sequential LP, etc.
	- 2nd order, sequential Newton
	- Lagrangian Relaxation, Primal-Dual, etc.

Other Classifications:

Quadratic, Convex, Integer, Mixed-Integer, Binary, etc.

Issues in Optimization

• **Problem Size**

- $-$ Small by hand
- Medium by software
- Large –by decomposition

• **Algorithm Complexity**

- Convergence speed
- Local Convergence speed

• **Insight more than just the solution**?

- Solution structure properties
- Sensitivity analysis
- Alternate formulations

What do you learn?

- Models the Art: intuition and common sense
	- How formulate real problems using quantitative models
- Theory the Science: theorems, geometries and universal rules
	- Necessary and Sufficient Conditions that must be true for the optimality of different classes of problems.
- Algorithms the Engineering: algorithms, methodologies and software tools
	- How we apply the theory to robustly and efficiently solve problems and gain insight beyond the solution.
- Applications AI, Machine Learning and Data Science
	- Logistic Regression, SVM, the Wasserstain barycenter, Reinforced learning/MDP, Information market,…

Art of Modeling, Formulation & Vocabulary

- Decision Variables $x \in \mathbb{R}^n$, yet to be decide
- Data/Coefficients, $c \in \mathbb{R}^n$, that are given and fixed
- Objective inner product $f = c^T x$: $\mathbb{R}^n \to \mathbb{R}$
- Constraint Set $X \subset \mathbb{R}^n$
- Feasible solution $x \in X$
- Optimal solution $x^* \in X^*$
- Optimal value $z^* = f(x^*)$

LP Example 1: Resource Allocation/Production Management

The Wyndor Glass Co. is a producer of high-quality glass products. It has three plants. Aluminum frames and hardware are made in Plant 1, wood frames are made in Plant 2, and Plant 3 is used to produce glass and assemble the products. Wyndor produces two products which require the resources of the three plants as follows:

max
$$
x_1 + 2x_2
$$

\ns.t. $x_1 \le 1$,
\n $2x_2 \le 2$,
\n $x_1 + x_2 \le 1.5$,
\n $x_1, x_2 \ge 0$

LP Example 2: Maximum Flow

How much flow can travel from A to B, given that each of the directed connecting routes have flow limits/capacities?

Maximum Flow by Inspection

How to certify that 15 is maximal?

Cut in Maximum Flow I

Cut value from Source site to Sink site=17

Cut in Maximum Flow II

Cut value from Source site to Sink site=17

Cut in Maximum Flow III

Cut value from Source site to Sink site=15

Data points classification application in Machine Learning and Data Science

Yinyu Ye, Stanford, MS&E211 Lecture Notes #1 17

Yinyu Ye, Stanford, MS&E211 Lecture Notes #1 18

LP Example 3: Transportation and Assignment

min
\n
$$
\sum_{i=1}^{3} \sum_{j=1}^{4} c_{ij} x_{ij}
$$
\n
$$
\sum_{j=1}^{4} x_{ij} = s_i, \quad \forall i = 1, 2, 3
$$
\n
$$
\sum_{i=1}^{3} x_{ij} = d_j, \quad \forall j = 1, 2, 3, 4
$$
\n
$$
x_{ij} \ge 0, \qquad \forall i, j
$$

Inventory Planning: *s* is part of the decision vars.

Machine Learning: The Wassestein Barycenter Problem I

The minimal transportation cost in Data Science is **WD(s, d^k)=** called the Wasserstein distance between a supply distribution and a demand distribution.

The **Wasserstein Barycenter Problem** is to find a distribution/points such that the sum of its Wasserstein distances to each of a set of distributions/points would be minimized

min**s ∑k WD(s, d^k) s.t. total mass constraint**

$$
\begin{aligned}\n\min \quad & \sum_{i=1}^{N} \sum_{j=1}^{N} c_{ij} x_{ij} \\
\text{s.t.} \quad & \sum_{j=1}^{N} x_{ij} = s_i, \quad \forall \ i = 1, \dots, N \\
& \sum_{i=1}^{N} x_{ij} = d_j, \quad \forall \ j = 1, \dots, N \\
& x_{ij} \geq 0, \quad \forall \ i, j\n\end{aligned}
$$

Machine Learning: The Wassestein Barycenter Problem II

What is the best ``mean or consensus'' image from a set of images (pixel distributions)?

- Simple average
- Simple average after re-centering
- The Wasserstein Barycenter of the set of images (self re-center and rotation)

Machine Learning: The Wassestein Barycenter Problem III

The simple avarage of n points is **s**=(∑_k **d**^k)/n or min_s ∑_{**k**} (|| **s – d^k||₂)²**

The **Wasserstein Barycenter Problem** is to find a distribution/points such that the sum of its Wasserstein distances to each of a set of distributions/points would be minimized (self re-center and rotation).

Simple average after re-centering Simple average the Barycenter image

LP Example 4: Support Vector Machine

x is the normal direction or slope vector and x_o is the intersect Find a line to **strictly** separate greens and reds

LP Example 4: Is Strict Separation Possible?

$$
a_i^T x + x_0 > 0, \forall i
$$

$$
b_j^T x + x_0 < 0, \forall j
$$

Are there x and x_0 such that the following (open) inequalities are all satisfied

$$
a_i^T x + x_0 \ge \varepsilon, \forall i
$$

$$
b_j^T x + x_0 \le -\varepsilon, \forall j
$$

Are there *x* and x_0 such that the following inequalities are all satisfied for arbitrarily small ε.

$$
\begin{array}{|l|l|}\n\hline\n\left(\frac{\partial f}{\partial x} + x_0 < 0, \forall j\right) & \text{(open) inequalities are all satisfied} \\
\hline\n\left(\frac{\partial f}{\partial y} + x_0 \geq \varepsilon, \forall i\right) & \text{Are there } x \text{ and } x_0 \text{ such that the following inequalities are all satisfied for arbitrarily small } \varepsilon.\n\end{array}
$$
\n
$$
\begin{array}{|l|}\n\hline\n\left(\frac{\partial f}{\partial y} + x_0 \leq -\varepsilon, \forall j\right) & \text{Divide } x \text{ and } x_0 \text{ by } \varepsilon, \text{ the problem can be equivalently reformulated.} \\
\hline\n\left(\frac{\partial f}{\partial y} + x_0 \leq -1, \forall j\right) & \text{Divide } x \text{ and } x_0 \text{ by } \varepsilon, \text{ the problem can be equivalently reformulated.} \\
\hline\n\left(\frac{\partial f}{\partial y} + x_0 \leq -1, \forall j\right) & \text{equivalently reformulated.} \\
\hline\n\left(\frac{\partial f}{\partial y} + x_0 \leq -1, \forall j\right) & \text{equivalently reformulated.} \\
\frac{\partial f}{\partial z} & \text{Standard, MSäE211 Lectures Notes #1}\n\end{array}
$$

Divide x and x_0 by ε ., the problem can be equivalently reformulated.

This is a special LP, called linear feasibility problem.

LP Example 4: Electric Vehicle Charging Schedule and Inventory Control

LP Example 5: When Discharge is Allowed

$$
\begin{array}{ll}\n\text{min} & \sum_{i=1}^{5} c_i x_i \\
\text{s.t.} & I_{i-1} + x_i - d_i = I_i, \quad \forall \ i = 1, 2, 3, 4, 5 \\
& I_{i-1} + x_i \le K, \quad \forall \ i = 1, 2, 3, 4, 5 \\
& I_i \ge 0, \quad \forall \ i.\n\end{array}
$$

Nonlinear Optimization: Bird or Plane?

Yinyu Ye, Stanford, MS&E211 Lecture Notes #1 27

Neural Network Design for Prediction

Back to Linear Programming

max
\n
$$
c_1x_1 + c_2x_2
$$
\n
$$
a_{11}x_1 + a_{12}x_2 \leq b_1
$$
\n
$$
a_{21}x_1 + a_{22}x_2 \leq b_2
$$
\n
$$
a_{31}x_1 + a_{32}x_2 \leq b_3
$$
\n
$$
x_1 \geq 0, x_2 \geq 0
$$

Abstract Linear Programming Model

max (min)
$$
c_1x_1 + c_2x_2 + ... + c_nx_n
$$

\ns.t. $a_{11}x_1 + a_{12}x_2 + ... + a_{1n}x_n \leq s = 0 \geq b_1$
\n $a_{21}x_1 + a_{22}x_2 + ... + a_{2n}x_n \leq s = 0 \geq b_2$
\n... ...

$$
a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \leq 0, \quad \sum_{n=1}^{\infty} a_n \geq 0, \quad \sum_{n=1}^{\infty} a_n \leq 0.
$$

Output : $x_1, ..., x_n$, decision v ariables a_{ii} , $i = 1,...,m$; $j = 1,...,n$, constraint left - hand - side table or matrix coef. Input: c_1 ,..., c_2 , objective coef.; b_1 ,..., b_m , constraint right - hand - side coef. *i j*

LP in Compact Matrix Form

s.t.
$$
A x {\leq, =, \geq} b
$$
,
 $x {\geq, \leq} 0 \text{ or free.}$

Some Facts of Linear Programming

- Add a constant to the objective function does not change the optimality
- Scale the objective coefficients does not change the optimality
- Scale the right-hand-side coefficients does not change the optimality but the solution scaled accordingly
- Reorder the decision variables (together with their corresponding objective and constraint coefficients) does not change the optimality
- Reorder the constraints (together with their right-hand-side coefficients) does not change the optimality
- Multiply both sides of an equality constraint by a constant does not change the optimality
- Pre-multiply both sides of all equality constraints by a non-singular matrix does not change the optimality