Optimal Transceiver Design for Multi-Access Communication

Lecturer: Tom Luo

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- An important problem in the management of communication networks: resource
- Frequency, transmitting power; Goal: high data rate, low bit error rate
- Transmitter + receiver for multi-user high speed broadband Digital Subscribe Line
- Direct formulation is nonconvex; equivalent formulation is SDP (thus convex);
- Further simplification to SOC, and to combinatorial polynomial time algorithm
- Valuable guidelines and insights for optimal practical transceiver design

Content

- Elements of data communication: OFDM, subcarriers, power loading, precoding/equalization.
- Linear transceiver (Transmitter + Receiver) design for the two user case:
- SDP formulation
- SOC formulation
- mdiinogle leimonylogly polynomial algorithm $O\left({{^8}n}
 ight)$
- General multi-user case
- Other formulations:
- Sum-rate transmitter design for multi-access channel
- Sum-rate transmitter design broadcast channel
- Linear transceiver design with zero-forcing equalizer

Single User SISO Communication System



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- ullet) f , g are transmitter filter and equalizer filter respectively
- əsion (neissue) əvitibbe əht zi n ullet

Material and the second second





- S/P: serial-parallel converter (with cyclic prefixing); P/S: parallel-serial
- $\mathbf{h} = \mathbf{h} \cdot \mathbf{h}$ transmitter matrix filter (or precoder), obtained from f; data rate $= \ell / n$. \mathbf{F}
- (artheta more than the term of term of
- \mathbf{H} : channel matrix (obtained from h); \mathbf{n} : noise
- $\mathbf{x} = \mathbf{H}^* \mathbf{F} \mathbf{s} + \mathbf{n}$, \mathbf{H}^* is circulant.



OFDM System

- The circulant channel matrix \mathbf{H}^* can be diagonalized via IFFT/FFT: $\mathbf{H}^* = \mathbf{D}^{\dagger}\mathbf{H}\mathbf{D}$, where \mathbf{D} is the standard discrete Fourier transform matrix, \mathbf{H} is diagonal
- The diagonalized channel becomes a set of independent subchannels
- Each subchannel corresponds to a subcarrier with a particular frequency (from FFT)
- This diagonalization is not channel dependent
- \bullet Orthogonal Frequency Division Multiplexing System employs IFFT/FFT to decompose the channel \mathbf{H}^*



Multi-Input Multi-Output Communication System

Linearly Precoded/Power Loaded OFDM

- The precoder F can be a general n imes n matrix, subject to power constraint ${
 m tr}({
 m FF}^{\dagger}) \leq p_{\rm L}$
- . The optimized design ${f F}$ will have a rank $\ell \geq \imath$, resulting in an optimal data rate ℓ / \imath .
- A special, and popular, linear precoder is the so called power loading precoder: F is diagonal.

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Linearly Precoded/Power Loaded OFDM





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Diagram of Two-user Communication System

 $\begin{array}{ll} \mbox{Mathematical model: } \mathbf{x} = \mathbf{H}_1 \mathbf{F}_1 \mathbf{s}_1 + \mathbf{H}_2 \mathbf{F}_2 \mathbf{s}_2 + \rho \mathbf{n}, & \rho > 0. \\ \mbox{Linear detection: } \mathbf{s}_i = \mathrm{sign} \left(\mathbf{G}_i \mathbf{x} \right), & i = 1, 2. \end{array}$

Goal: Given the channel matrices, \mathbf{H}_1 , \mathbf{H}_2 , design transceivers \mathbf{F}_1 , \mathbf{F}_2 , \mathbf{G}_1 , \mathbf{G}_2 .

Applications and Previous Work

- Applications include the current and future systems of DSL, DAB, DVB.
- Equalizer (or receiver) design for a fixed transmitter has been researched extensively in the last decade.
- The joint transmitter and receiver (transceiver) design was considered recently, but only
- In the single user transceiver design work, the design criteria used include:
- Minimum Mean Square Error,
- , maximum information rate,
- Vticeqec lennedo –
- The last two require complex receiver structures.

Mean Square Error

- . ($\mathbf{x}_i\mathbf{B}$) $\mathrm{ngis}={}_i\hat{\mathbf{s}}$: ${}_i\mathbf{D}$ (receiver (equalizer) denotes the equation of the equat
- Let e_i denote the error vector (before making the hard decision) for user i, i = 1, 2. Then

$$\mathbf{e}_{1} = \mathbf{G}_{1}\mathbf{x} - \mathbf{s}_{1} = \mathbf{G}_{1}(\mathbf{H}_{1}\mathbf{F}_{1}\mathbf{s}_{1} + \mathbf{H}_{2}\mathbf{F}_{2}\mathbf{s}_{2} + \rho\mathbf{G}_{1}\mathbf{n}.$$
$$= (\mathbf{G}_{1}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I})\mathbf{s}_{1} + \mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}\mathbf{s}_{2} + \rho\mathbf{G}_{1}\mathbf{n}.$$

• This further implies

$$E(\mathbf{e}_{1}\mathbf{e}_{1}^{1}) = (\mathbf{G}_{1}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I}) \left(\mathbf{G}_{1}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I}\right)_{\dagger} + \left(\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}\right) \left(\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}\right)_{\dagger} + \rho^{2}\mathbf{G}_{1}\mathbf{G}_{\dagger}^{1}$$

• Similarly, we have

$$E(\mathbf{e}_{2}\mathbf{e}_{2}^{2}) = (\mathbf{G}_{2}\mathbf{H}_{2}\mathbf{F}_{2} - \mathbf{I}) \left(\mathbf{G}_{2}\mathbf{H}_{2}\mathbf{F}_{2} - \mathbf{I}\right)^{\dagger} + \left(\mathbf{G}_{2}\mathbf{H}_{1}\mathbf{F}_{1}\right) \left(\mathbf{G}_{2}\mathbf{H}_{1}\mathbf{F}_{1}\right)^{\dagger} + \rho^{2}\mathbf{G}_{2}\mathbf{G}_{2}^{2}.$$

D.Q. Luo

Formulation: MMSE Equalizer Case

• Our goal is to design a set of transmitting matrix filters ${f F}_i$ and a set of matrix equalizers ${f G}_i$ such that the total mean squared error

$$MSE = tr(E(\mathbf{e}_1\mathbf{e}_1^{\dagger})) + tr(E(\mathbf{e}_2\mathbf{e}_2^{\dagger}))$$

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 As is always the case in practice, there are power constraints on the transmitting matrix filters:

$$\operatorname{tr}(\mathbf{F}_1\mathbf{F}_1^{\intercal}) \leq p_1, \quad \operatorname{tr}(\mathbf{F}_2\mathbf{F}_2^{\intercal}) \leq p_2$$

- The above is nonconvex.
- \bullet We first eliminate the variables $\mathbf{G}_{1,}$ $\mathbf{G}_{2;}$ the MMSE equalizers.

Formulation: MMSE Equalizer Case

• By minimizing $E(e_1e_1^{\dagger})$ with respect to G_1 , we obtain the following MMSE equalizer for user 1: $G_1 = F_1^{\dagger}H_1^{\dagger}W$, where

$$\mathbf{W} = \left(\mathbf{H}_{1}\mathbf{F}_{1}\mathbf{F}_{1}^{\dagger}\mathbf{H}_{1}^{\dagger}\mathbf{H}_{2}^{\dagger}\mathbf{H}_{2}\mathbf{F}_{2}\mathbf{F}_{2}\mathbf{F}_{2}\mathbf{F}_{2}^{\dagger}\mathbf{H}_{2}^{\dagger}\mathbf{H}_{2}^{\dagger}\right)^{-1}.$$

- Substituting this into $E(\mathbf{e}_1\mathbf{e}_1^\dagger)$ gives:

$$E(\mathbf{e}_{1}\mathbf{e}_{1}^{\dagger}) = -\mathbf{F}_{1}^{\dagger}\mathbf{H}_{1}^{\dagger}\mathbf{M}\mathbf{H}_{1}\mathbf{I}\mathbf{I} = -\mathbf{I}_{1}^{\dagger}\mathbf{e}_{1}\mathbf{H}_{1}$$

• Similarly, the MMSE equalizer G_2 for user 2 is given by $G_2 = F_2^{\dagger}H_2^{\dagger}N$ and resulting

$$E(\mathbf{e}_2\mathbf{e}_2^{\dagger}) = -\mathbf{F}_2^{\dagger}\mathbf{H}_2^{\dagger}\mathbf{W}\mathbf{H}_2\mathbf{F}_2 + \mathbf{I}.$$

32M Interview

Substituting into the above expression gives rise to

MSE =
$$\operatorname{tr}(E(\mathbf{e}_{1}\mathbf{e}_{1}^{\dagger})) + \operatorname{tr}(E(\mathbf{e}_{2}\mathbf{e}_{2}^{2})) + \operatorname{tr}(E(\mathbf{e}_{2}\mathbf{e}_{2}^{2})) + 2n$$

= $-\operatorname{tr}\left(\mathbf{W}(\mathbf{H}_{1}\mathbf{F}_{1}\mathbf{F}_{1}^{\dagger}\mathbf{H}_{1}^{\dagger}) - \operatorname{tr}\left(\mathbf{W}\mathbf{H}_{2}\mathbf{F}_{2}\mathbf{F}_{2}^{\dagger}\mathbf{H}_{2}^{\dagger}\right) + 2n$
= $-\operatorname{tr}\left(\mathbf{W}(\mathbf{H}_{1}\mathbf{F}_{1}\mathbf{F}_{1}^{\dagger}\mathbf{H}_{1}^{\dagger}) - \operatorname{tr}\left(\mathbf{W}\mathbf{H}_{2}\mathbf{F}_{2}\mathbf{F}_{2}^{\dagger}\mathbf{H}_{2}^{\dagger}\right) + 2n$

where the last step follows from the definition of ${f W}.$

Formulation: MMSE Equalizer Case

- Introduce matrix variables: $\mathbf{U}_1 = \mathbf{F}_1 \mathbf{F}_1^{\dagger}$, $\mathbf{U}_2 = \mathbf{F}_2 \mathbf{F}_2^{\dagger}$.
- Then the MMSE transceiver design problem becomes

 $\begin{array}{ll} \text{minimize}_{\mathbf{U}_1,\mathbf{U}_2} & \text{tr}\left((\mathbf{H}_1\mathbf{U}_1\mathbf{H}_1^{\dagger} + \mathbf{H}_2\mathbf{U}_2\mathbf{H}_2^{\dagger} + \rho^2\mathbf{I})^{-1}\right) \\ \text{subject to} & \text{tr}(\mathbf{U}_1) \leq p_1, \quad \text{tr}(\mathbf{U}_2) \leq p_2, \\ \mathbf{U}_1 \geq \mathbf{0}, \quad \mathbf{U}_2 \geq \mathbf{0}. \end{array}$

ullet Beformulate using the auxiliary matrix variable ${f W}$:

 $\begin{array}{ll} \text{minimize} W, \textbf{U}_1, \textbf{U}_2 & \text{tr} (\mathbf{W}) \\ \text{subject to} & \text{tr} (\mathbf{U}_1) \leq p_1, \quad \text{tr} (\mathbf{U}_2) \leq p_2, \\ \mathbf{W} \succeq (\mathbf{H}_1 \mathbf{U}_1 \mathbf{H}_1^{\dagger} + \mathbf{H}_2 \mathbf{U}_2 \mathbf{H}_2^{\dagger} + \rho^2 \mathbf{I})^{-1} \\ \mathbf{U}_1 \succeq \mathbf{0}, \quad \mathbf{U}_2 \succeq \mathbf{0}. \end{array}$

• The constraint $\mathbf{W} \succeq \mathbf{W} \succeq \mathbf{W}_1 \mathbf{U}_1 \mathbf{H}_1 \mathbf{U}_2 \mathbf{U}_2 \mathbf{H} + \frac{\dagger}{2} \mathbf{H}_2 \mathbf{U}_2 \mathbf{H} + \frac{\dagger}{2} \mathbf{H}_1 \mathbf{U}_1 \mathbf{H}) \preceq \mathbf{W}$ in equivalent to LMI:

(E)
$$\mathbf{I} = \mathbf{I}_{\mathbf{I}} \mathbf{U}_{\mathbf{I}} \mathbf{H}_{\mathbf{I}} \mathbf{U}_{\mathbf{I}} \mathbf{H}_{\mathbf{I}}^{\dagger} + \mathbf{H}_{\mathbf{I}} \mathbf{U}_{\mathbf{I}} \mathbf{H}_{\mathbf{I}}^{\dagger} + \mathbf{\rho}^{2} \mathbf{I}$$
 (3)

• We obtain an SDP formulation:

 $\begin{array}{lll} \text{minimize} W, U_1, U_2 & \text{tr} (W) \\ \text{subject to} & \text{tr} (U_1) \leq p_1, & \text{tr} (U_2) \leq p_2, \\ W & \text{satisfies (3)}, \\ U_1 \succeq 0, & U_2 \succeq 0. \end{array}$

• Interior point method with arithmetic complexity $O(n^{6.5}\log(1/\epsilon)),\ \epsilon>0$ is the solution accuracy.

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Result

If \mathbf{H}_1 and \mathbf{H}_2 are diagonal, as in the OFDM systems, then the optimal transmitters are also diagonal.

Implication

The MMSE transceivers for an multi-user OFDM system can be implemented by optimally setting the data rates and allocating power to each subcarrier for all the users.

Linearly Precoded/Power Loaded OFDM





From SDP to SOC Formulation

Restricting to diagonal designs, the SDP becomes SOC:

$$\begin{array}{ll} \text{minimize}_{\mathbf{w},\mathbf{u}_{1},\mathbf{u}_{2}} & \sum_{i=1}^{n} \mathbf{w}_{i} \\ \text{subject to} & \sum_{i=1}^{n} \mathbf{u}_{1}(i) \leq p_{1}, & \sum_{i=1}^{n} \mathbf{u}_{2}(i) \leq p_{2}, \\ \mathbf{w}_{i}\left(|\mathbf{h}_{1}(i)|^{2}\mathbf{u}_{1}(i) + \mathbf{h}_{2}(i)|^{2}\mathbf{u}_{2}(i) + \mathbf{p}^{2}\right) \geq 1, \\ \mathbf{u}_{1}(i) \geq 0, & \mathbf{u}_{2}(i) \geq 0, & i = 1, 2, ..., n. \end{array}$$

- There exist highly efficient (general purpose) interior point methods to solve the above second order cone program.
- Arithmetic complexity $O(n^{3.5}\log(1/\epsilon))$, $\epsilon>0$ is the accuracy.

Properties of Optimal MMSE Transceiver

• Let $u_1^* \ge 0, \ u_2^* \ge 0$ be the optimal transceivers. Define:

$$\begin{cases} I_{1} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{2} = \{i \mid \mathbf{u}_{1}^{*}(i) = 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{3} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{4} = \{i \mid \mathbf{u}_{1}^{*}(i) = 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{5} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{6} = \{i \mid \mathbf{u}_{1}^{*}(i) = 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{6} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{6} = \{i \mid \mathbf{u}_{1}^{*}(i) = 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{6} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{6} = \{i \mid \mathbf{u}_{1}^{*}(i) = 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \\ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{2}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7} = \{i \mid \mathbf{u}_{1}^{*}(i) > 0, \ \mathbf{u}_{1}^{*}(i) > 0 \}, \ I_{7}$$

- I_1 , I_2 : subcarriers allocated to user 1 and user 2; I_s and I_u : subcarriers *shared* and *unused*; data rates: $(|I_1| + |I_s|)/n$, $(|I_2| + |I_s|)/n$
- For each $i \in I_1$ and $j \in I_2$, we have $\frac{|\mathbf{h}_1(i)|^2}{|\mathbf{h}_2(i)|^2} \ge \frac{|\mathbf{h}_2(j)|^2}{|\mathbf{h}_2(j)|^2}$.
- For any $i \in I_u$ and any $j \in I_1 \cup I_s$, we have $|\mathbf{h}_1(i)|^2 < |\mathbf{h}_1(j)|^2$. Similarly, for any $i \in I_u$ and any $j \in I_2 \cup I_s$, we have $|\mathbf{h}_2(i)|^2 < |\mathbf{h}_2(j)|^2$.

Intuitive Interpretation

- $\mathbf{h}_2(i)\mathbf{f}_2(i)\mathbf{s}_2(i) + \mathbf{P}_2\mathbf{F}_2\mathbf{s}_2 + \rho\mathbf{n}$, with \mathbf{H}_i , \mathbf{F}_i diagonal; $\mathbf{x}(i) = \mathbf{h}_1(i)\mathbf{f}_1(i)\mathbf{s}_1(i) + \mathbf{h}_2\mathbf{h}_2\mathbf{n}(i)$.
- In a fading environment, the path gains $|\mathbf{h}_1(i)|^2$, $|\mathbf{h}_2(i)|^2$ are random,
- ⇒ the probability of having two equal path gains is zero.
- $rac{1}{s}$ is singleton: at most one subcarrier should be shared by the two users.
- The remaining subcarriers are allocated to the two users according to the path gain ratios: subcarrier i to user 1 and subcarrier j to user 2 only if

$$\frac{|\mathbf{p}^{z}(i)|_{z}}{|\mathbf{p}^{z}(i)|_{z}} \ge \frac{|\mathbf{p}^{z}(i)|_{z}}{|\mathbf{p}^{z}(i)|_{z}}$$

• The subcarriers in I_u have small path gains for both users (i.e., both $|\mathbf{h}_1(i)|^2$ and $|\mathbf{h}_2(i)|^2$ are subcarriers! subcarriers!

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- The properties of optimal MMSE transceivers can be used to design a combinatorial algorithm.
- Asymptotic equation is a set of the equation of the equation is a set of the equation is a se
- .i amos not $\{n\,,\ldots,i\}\supseteq I_2\subseteq \{i\,,\ldots,1\}\supseteq I_1$ nad T ullet
- Leads to an $O(n^3)$ strongly polynomial time (combinatorial) algorithm (vs. $O(n^{3.5}\log 1/\epsilon)$ interior point algorithm for SOC).

Practical Implications



General m-User Case

Mathematical model:

$$\mathbf{x} = \mathbf{H}_{1}\mathbf{F}_{1}\mathbf{s}_{1} + \mathbf{H}_{2}\mathbf{F}_{2}\mathbf{s}_{2} + \dots + \mathbf{H}_{m}\mathbf{F}_{m}\mathbf{s}_{m} + \rho\mathbf{n}.$$

Let G_i be the linear MMSE matrix equalizer at the *i*-th receiver. Then the total MSE is given by

$$\rho^{2} \operatorname{tr} \left(\mathrm{I} - m \right) + \left(\mathrm{I}^{-} (\mathbf{I}^{2} \varphi + \mathrm{I}^{\dagger} \mathbf{H}_{i}^{\dagger} \mathbf{H}_{i}^{\dagger} \mathbf{H}_{i}^{\dagger} \mathbf{H}_{i} \mathbf{H} + \dots + \mathrm{I}^{\dagger} \mathbf{H}_{1}^{\dagger} \mathbf{H}_{1}^{\dagger} \mathbf{H}_{1} \mathbf{H} \right) \right) \operatorname{tr}^{2} \varphi$$

• Let $\mathbf{U}_i=\mathbf{F}_i\mathbf{F}_i^\dagger$. Then the power constrained optimal MMSE transmitter design

$$\begin{array}{ll} \text{minimize} U^{1-}(\mathbf{I}^2\boldsymbol{\eta} + \overset{\dagger}{m} \mathbf{H}_m \mathbf{U}_m \mathbf{H} + \cdots + \overset{\dagger}{\mathbf{I}} \mathbf{H}_1 \mathbf{U}_1 \mathbf{H}) & \text{tr} & (\mathbf{H}^2, \cdots, \mathbf{U}_m) \\ \text{minimize} \mathbf{U}_i, \cdots, \mathbf{I} & = i & \mathbf{U}_i \succeq \mathbf{0}, \quad \mathbf{U}_i \succeq \mathbf{0}, \quad \mathbf{i} = 1, \dots, m. \end{array}$$

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SDP formulation

minimize tr
$$(\mathbf{W})$$

subject to $\operatorname{tr}(\mathbf{U}_i) \leq p_i, \quad \mathbf{U}_i \succeq \mathbf{0}, \quad i = 1, 2, ..., m,$
 \mathbf{W}
 \mathbf{W}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}
 $\mathbf{H}_m \mathbf{U}_m \mathbf{H} + \dots + \frac{\dagger}{I} \mathbf{H}_1 \mathbf{U}_1 \mathbf{H}$
 \mathbf{I}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}
 \mathbf{M}

SOC formulation

Simulation Scenario

- Uplink with 16 active users and 160 available subcarriers
- Each user "sees" its own Rayleigh channel (complex-valued)
- Three schemes:
- AMOUR No channel knowledge; Each user uses 10 subcarriers, spreads 8 bits over these carriers using a DFT-type spreading.
- Individually MMSE power-loaded OFDM Same subcarrier allocation as AMOUR.
 Each user sends 1 bit per subcarrier, i.e 10 bits per block; knows its allocated channels and does MMSE power loading for these bits.
- Multi-user MMSE power loaded OFDM Using the SOCP formulation. In this case the subcarrier allocations and the number of bits per block vary from block to block, but the average number of bits per block remains 10.

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Efficiency of the Design Approach

, On a PIII 600Mhz PC,

- Two users, 2 symbols per block, length 3 channel;
- SDP \sim 0.65 secs
- SOCP \sim 0.13 secs
- 16 users, 10 symbols per block, length 3 channel
- sces = 0.65 \sim 9.05 -

Formulation: Max Sum Rate Capacity for MAC

- Let $\Sigma_k \succeq 0$, p_k denote the covariance matrix and the transmit power of the k-th user signal.
- The total sum rate of multi-access channel is

$$({}^{\mathcal{H}}_{\mathcal{H}}\mathbf{H}_{\mathcal{H}}\mathbf{Z}_{\mathcal{H}}\mathbf{H}\sum_{\mathbf{I}=\mathcal{H}}^{\mathcal{H}}\mathbf{H}(\mathbf{I}+\mathbf{I})$$
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which is achievable by successive nulling and cancellation at BS.

• The multi-user transmitter design is then

maximize log det $(\mathbf{I} + \sum_{k=1}^{K} \mathbf{H}_{k} \mathbf{\Sigma}_{k} \mathbf{H}_{k} \mathbf{L}_{k} \mathbf{I}_{k})$ and maximize log det $(\mathbf{I} + \sum_{k=1}^{K} \mathbf{H}_{k} \mathbf{\Sigma}_{k} \mathbf{Z}_{k})$ and $\mathbf{I} = \mathbf{1}, \mathbf{2}, \dots, K$.

• A convex problem; can be solved by interior point methods, iterative water-filling.

Formulation: Max Sum Rate Capacity for BC

- Let $\Sigma_k \succeq 0$ denote the covariance matrix for the k-th user signal, and let p denote the total transmit power.
- By duality, total sum rate of a broadcast channel is

$$\log \det(\mathbf{I} + \sum_{k=1}^{K} \mathbf{H}_{k}^{H} \bar{\mathbf{\Sigma}}_{k}^{H} \mathbf{H}_{k})$$

which is achievable by dirty paper coding technique, where Σ_k (new variables) depends on Σ_k linearly.

The multi-user transmitter design is then

maximize log det $(\mathbf{I} + \sum_{k=1}^{K} \mathbf{H}_{k} \overline{\boldsymbol{\Sigma}}_{k} \mathbf{H}_{k} \overline{\boldsymbol{\Sigma}}_{k} \mathbf{H}_{k}^{H})$ subject to $\sum_{k=1}^{K} \operatorname{tr}(\overline{\boldsymbol{\Sigma}}_{k}) \leq p, \quad \overline{\boldsymbol{\Sigma}}_{k} \geq 0, \quad k = 1, 2, ..., K.$

• A convex problem; can be solved by interior point methods (and iterative water-filling?).

Formulation: Zero-Forcing Equalizer Case

• Recall e_i denotes the error vector (before making the hard decision) for user i, i = 1, 2and

$$\mathbf{e}_1 = \mathbf{G}_1 \mathbf{x} - \mathbf{s}_1 = \mathbf{G}_1 (\mathbf{H}_1 \mathbf{F}_1 \mathbf{s}_1 + \mathbf{H}_2 \mathbf{F}_2 \mathbf{s}_2 + \rho \mathbf{G}_1 \mathbf{n}.$$

= $(\mathbf{G}_1 \mathbf{H}_1 \mathbf{F}_1 - \mathbf{I}) \mathbf{s}_1 + \mathbf{G}_1 \mathbf{H}_2 \mathbf{F}_2 \mathbf{s}_2 + \rho \mathbf{G}_1 \mathbf{n}.$

Moreover,

$$E(\mathbf{e}_{1}\mathbf{e}_{1}^{\dagger}) = (\mathbf{G}_{1}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I}) (\mathbf{G}_{1}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I})^{\dagger} + (\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}) (\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2})^{\dagger} + \rho^{2}\mathbf{G}_{2}\mathbf{G}_{1}^{\dagger}.$$

$$E(\mathbf{e}_{2}\mathbf{e}_{1}^{\dagger}) = (\mathbf{G}_{2}\mathbf{H}_{2}\mathbf{F}_{2} - \mathbf{I}) (\mathbf{G}_{2}\mathbf{H}_{1}\mathbf{F}_{1} - \mathbf{I})^{\dagger} + (\mathbf{G}_{2}\mathbf{H}_{1}\mathbf{F}_{2}) (\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2})^{\dagger} + \rho^{2}\mathbf{G}_{2}\mathbf{G}_{1}^{\dagger}.$$

$$\mathbf{G}_1 = (\mathbf{H}_1\mathbf{F}_1)^{-1}, \quad \mathbf{G}_2 = (\mathbf{H}_2\mathbf{F}_2)^{-1}.$$

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Formulation: Zero-Forcing Equalizer Case

Substituting the ZF conditions into the MSE expressions gives

$$MSE = \operatorname{tr} \left((\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}) (\mathbf{G}_{1}\mathbf{H}_{2}\mathbf{F}_{2}) \left(\mathbf{G}_{2}\mathbf{H}_{1}\mathbf{F}_{1} \right)_{\dagger} \right) + \rho^{2} \operatorname{tr} \left(\mathbf{G}_{2}\mathbf{G}_{\dagger}^{1} \right) + h^{2} \operatorname{tr} \left(\mathbf{G}_{2}\mathbf{G}_{\dagger}^{2} \right) + h^{2} \operatorname{$$

Introduce new matrix variables

$$\mathbf{U}_{1} = \mathbf{F}_{1}\mathbf{F}_{\dagger}^{1}, \quad \mathbf{U}_{2} = \mathbf{F}_{2}\mathbf{F}_{\dagger}^{2}, \quad \mathbf{V}_{1} = \mathbf{G}_{\dagger}^{\dagger}\mathbf{G}_{1}, \quad \mathbf{V}_{2} = \mathbf{G}_{\dagger}^{2}\mathbf{G}_{2}.$$

• Then the MSE can be rewritten as

 $\mathsf{MSE} = \mathsf{tr}(\mathbf{V}_1\mathbf{H}_2\mathbf{U}_2\mathbf{H}_2^{\dagger}) + \rho^2 \mathsf{tr}(\mathbf{V}_1) + \mathsf{tr}(\mathbf{V}_2\mathbf{H}_1\mathbf{U}_1\mathbf{H}_1^{\dagger}) + \rho^2 \mathsf{tr}(\mathbf{V}_2)$

- The power constraint becomes $tr(U_1) \leq p_1$, $tr(U_2) \leq p_2$.
- The ZF condition reduces to $\mathbf{V}_1^{-1} = \mathbf{H}_1 \mathbf{U}_1 \mathbf{H}_1^{\dagger}$, $\mathbf{V}_2^{-1} = \mathbf{H}_2 \mathbf{U}_2 \mathbf{H}_2^{\dagger}$.

Formulation: Zero-Forcing Case

• The Minimum MSE transceiver design problem can be cast as

| $\mathbf{V}_i \succeq 0, \mathbf{U}_i \succeq 0, i = 1, 2.$ | |
|---|------------|
| $\mathbf{V}_1^{\mathrm{I}} = \mathbf{H}_1 \mathbf{U}_1 \mathbf{H}_{1}^{\mathrm{I}}, \mathbf{V}_2^{\mathrm{I}} = \mathbf{H}_2 \mathbf{U}_2 \mathbf{H}_{2}^{\mathrm{I}},$ | |
| $\operatorname{tr}(\mathrm{U}_1)\leq p_1,\operatorname{tr}(\mathrm{U}_2)\leq p_2,$ | subject to |
| $MSE = \operatorname{tr}(\mathbf{V}_{1}\mathbf{H}_{2}\mathbf{U}_{2}\mathbf{H}_{2}^{\dagger}) + \rho^{2}\operatorname{tr}(\mathbf{V}_{1}) + \operatorname{tr}(\mathbf{V}_{2}\mathbf{H}_{1}\mathbf{U}_{1}\mathbf{H}_{1}^{\dagger}) + \rho^{2}\operatorname{tr}(\mathbf{V}_{2})$ | əziminim |

- Note the constraints are nonlinear (due to the matrix inverse)
- The objective function is nonconvex quadratic, due to the cross terms $tr(\mathbf{V}_1\mathbf{H}_2\mathbf{U}_2\mathbf{H}_2\mathbf{H}_2)$.
- Reformulation is necessary.

Reformulation: ZF Case

 Use monotonicity and Schur complement technique, we obtain the following equivalent formulation:

$$\begin{array}{ll} \text{minimize} & \text{MSE} = \operatorname{tr}(\mathbf{V}_1\mathbf{H}_2\mathbf{U}_2\mathbf{H}_2^1) + \rho^2\operatorname{tr}(\mathbf{V}_1) + \operatorname{tr}(\mathbf{V}_2\mathbf{H}_1\mathbf{U}_1\mathbf{H}_1^1) + \rho^2\operatorname{tr}(\mathbf{V}_2) \\ \text{subject to} & \operatorname{tr}(\mathbf{U}_1) \leq p_1, \quad \operatorname{tr}(\mathbf{U}_2) \leq p_2, \quad \mathbf{V}_i \succeq \mathbf{0}, \quad \mathbf{U}_i \succeq \mathbf{0}, \quad i = 1, 2 \\ & \left[\begin{array}{cc} \mathbf{H}_1\mathbf{U}_1\mathbf{H}_1^{\dagger} & \mathbf{I} \\ \mathbf{I} & \mathbf{V}_1 \end{array} \right] \succeq \mathbf{0}, \\ & \left[\begin{array}{cc} \mathbf{H}_2\mathbf{U}_2\mathbf{H}_2^{\dagger} & \mathbf{I} \\ \mathbf{I} & \mathbf{V}_2 \end{array} \right] \succeq \mathbf{0}. \end{array} \right.$$

- Note that the constraints are all linear matrix inequalities (LMIs), and in particular convex.
- But the objective function is nonconvex.

Alternating Direction Method

- Fixing the designs for user 1 (namely, U_1 , V_1), the objective function MSE is linear in U_2 , V_2 , resulting in a SDP. Similarly, fixing U_2 and V_2 yields a semidefinite program in U_1 , U_1 , and V_1 .
- Alternating Direction Method:

,
$$\mathbf{I} \leq \lambda$$
 noiseration $\mathbf{I} = \mathbf{V}_i^{(0)} = \mathbf{V}_i^{(0)} = \mathbf{I}$. At iteration $k \geq 1$,

- * Solve (1) with U_2 , V_2 fixed to the values of $U_2^{(k-1)}$, $V_2^{(k-1)}$. Update U_1 and V_1 . $V_1^{(k)}$ to the resulting optimized values of U_1 and V_1 .
- * Solve (1) with U_1 , V_1 fixed to the values of $U_1^{(k-1)}$, $V_1^{(k-1)}$. Update $U_2^{(k)}$ and $V_2^{(k)}$ to the resulting optimized values of U_2 and V_2 .
- . I + λ =: λ hith the sequence of the seq
- Convergence: bounded iterates + the minimum principle necessary optimality condition.

Power-loaded OFDM Optimal?

- Let channel matrices $\mathbf{H}_1, \mathbf{H}_2$ be diagonal.
- If we fix \mathbf{U}_2 , \mathbf{V}_2 at some positive definite diagonal matrices in (1) and optimize with respect to \mathbf{U}_1 , \mathbf{V}_1 , then the resulting optimized matrices \mathbf{U}_1 , \mathbf{V}_1 , can also be taken to be positive definite and diagonal.
- The proof uses reduction and a property of bipartite matching polytope.
- Conjecture: the optimal solutions of (1) are always diagonal.
- Imply the power-loaded OFDM is optimal.

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- (lenogeib ${}_i \mathbf{V}$, ${}_i \mathbf{U}$) angiseb lenogeib of tritter (U \bullet
- The formulation (1) reduces to a geometric program:

minimize
$$\sum_{\substack{i=1\\i=1}}^{n} \left(\mathbf{v}_{1}^{-1}(i)\mathbf{v}_{2}(i) + \mathbf{v}_{1}(i)\mathbf{v}_{2}^{-1}(i)\right) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \mathbf{v}_{2}(i)) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i)) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i)) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i) + \rho^{2} \sum_{i=1}^{n} (\mathbf{v}_{1}(i)) + \rho^{2} \sum_{i=1}^{n} (\mathbf$$

- (2) can be turned into a convex program by using the standard logarithmic transformation.
- The dual of (2) is a linearly constrained entropy maximization problem.

The Dual Program

maximize
$$-\sum_{i=1}^{6n} \delta_i \log \delta_i - \sum_{i=1}^{2n} c_i \delta_{4n+i} + \lambda_1 \log \lambda_1 + \lambda_2 \log \lambda_2$$

subject to
$$\delta_{2i-1} - \delta_{2i} + \rho^2 \delta_{2n+2i-1} - \delta_{4n+i} = 0, \quad 1 \leq i \leq n,$$
$$\lambda_1 = \sum_{i=4n+1}^{5n} \delta_i, \quad \lambda_2 = \sum_{i=5n+1}^{6n} \delta_i,$$
$$\lambda_1 = \sum_{i=4n+1}^{5n} \delta_i, \quad \lambda_2 = \sum_{i=5n+1}^{6n} \delta_i,$$
$$\sum_{i=1}^{4n} \delta_i = 1, \quad \delta_i \geq 0, \quad 1 \leq i \leq 6n,$$

where the coefficients c_i are defined as

$$c_i = \left\{ \begin{array}{ccc} & i \geq 1 & , & ((i)_1^2 \mathbf{h}_1 q) \text{ gol} \\ & \log \left(p_2 \mathbf{h}_2^2 (i) \right) & , & 1 \leq i \leq n, \\ & \log \left(p_2 \mathbf{h}_2^2 \mathbf{h}_2 q \right) \text{ gol} \end{array} \right\} = iO$$

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- So far we have
- Presented various SDP/SOC formulations and algorithms for the optimal transceiver
- Studied the properties of the optimal transceiver designs.
- Demonstrated the potential of SDP/SOC/interior point methods in digital communication.
- Results provide valuable guidelines and insights for the practical system design.
- Future work
- Incorporating QoS and other receiver structures in the formulation.
- Extension to the multi-user downlink case.