# Programming Abstractions <br> CS106B 

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## Graphs Topics

Graphs!

1. Basics

- What are they? How do we represent them?

2. Theorems

- What are some things we can prove about graphs?

3. Breadth-first search on a graph

- Spoiler: just a very, very small change to tree version

4. Dijkstra's shortest paths algorithm

- Spoiler: just a very, very small change to BFS

5. A* shortest paths algorithm (continued)

- Spoiler: just a very, very small change to Dijkstra's

6. Minimum Spanning Tree

- Kruskal's algorithm


## A* Solving Super Mario (video)

https://youtu.be/DIkMs4ZHHr8

What goes in the ??? ?
A. $2+5$ ?
B. $1+6$ ?
C. $2+4$ ?
D. Other/none/more

$A^{*}$ : enqueue neighbors.



Now we're done with the green " 1 " node's turn.

What is the next node to turn green? (and what would it be if this were Dijkstra's?)
$A^{*}$ : dequeue next lowest priority value node. Notice we are making a straight line right for the end point, not wasting time with other directions.
$A^{*}$ : enqueue neighbors-uh-oh, wall blocks us from continuing forward.

$A^{*}$ : eventually figures out how to go around the wall, with some waste in each direction.

|  |  |  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | $\begin{aligned} & 4+ \\ & 7 ? \end{aligned}$ | $\begin{aligned} & 5+ \\ & 6 ? \end{aligned}$ | $\begin{aligned} & 6+ \\ & 5 ? \end{aligned}$ | $\begin{aligned} & 7+ \\ & 4 ? \end{aligned}$ |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | 2 | 3 | 4 | 5 | 6 | $\begin{aligned} & 7+ \\ & 2 ? \end{aligned}$ |  |
|  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | 2 | 1 | 2 | 3 |  | $\begin{aligned} & 7+ \\ & 2 ? \end{aligned}$ |  |  |
| $3+$ | 2 | 1 | 2 | 1 | 2 |  | 8 | 2 |  |
|  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | 2 | 1 | 2 | 3 |  | 7 | $\begin{aligned} & 8+ \\ & 1 ? \end{aligned}$ |  |
|  |  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | 2 | 3 | 4 | 5 | 6 | 7 | $\begin{aligned} & 8+ \\ & 3 ? \end{aligned}$ |
|  |  |  | $\begin{aligned} & 3+ \\ & 8 ? \end{aligned}$ | $4+$ | $\begin{aligned} & 5+ \\ & 6 ? \end{aligned}$ | $\begin{aligned} & 6+ \\ & 5 ? \end{aligned}$ | $\begin{aligned} & 7+ \\ & 4 ? \end{aligned}$ | $\begin{aligned} & 8+ \\ & 3 ? \end{aligned}$ |  |

For Comparison: What Dijkstra's Algorithm Would Have Searched


- Mark all nodes as gray.
- Mark the initial node s as yellow and at candidate distance 0.
- Enqueue s into the priority queue with priority 0.


## Dijkstra's Algorithm

- While not all nodes have been visited:
- Dequeue the lowest-cost node $u$ from the priority queue.
- Color $u$ green. The candidate distance $d$ that is currently stored for node $u$ is the length of the shortest path from $s$ to $u$.
- If $u$ is the destination node $t$, you have found the shortest path from $s$ to $t$ and are done.
- For each node $v$ connected to $u$ by an edge of length $L$ :
- If $v$ is gray:
- Color $v$ yellow.
- Mark v's distance as $d+L$.
- Set $v$ 's parent to be $u$.
- Enqueue $v$ into the priority queue with priority $d+L$.
- If $v$ is yellow and the candidate distance to $v$ is greater than $d+L$ :
- Update $v$ 's candidate distance to be $d+L$.
- Update v's parent to be $u$.
- Update $v$ 's priority in the priority queue to $d+L$.
- Mark all nodes as gray.
- Mark the initial node $s$ as yellow and at candidate distance 0.


## A* Search

- Enqueue s into the priority queue with priority h(s,t).
- While not all nodes have been visited:
- Dequeue the lowest-cost node $u$ from the priority queue.
- Color $u$ green. The candidate distance $d$ that is currently stored for node $u$ is the length of the shortest path from $s$ to $u$.
- If $u$ is the destination node $t$, you have found the shortest path from $s$ to $t$ and are done.
- For each node $v$ connected to $u$ by an edge of length $L$ :
- If $v$ is gray:
- Color v yellow.
- Mark v's distance as $d+L$.
- Set $v$ 's parent to be $u$.
- Enqueue $v$ into the priority queue with priority $d+L+h(v, t)$.
- If $v$ is yellow and the candidate distance to $v$ is greater than $d+L$ :
- Update $v$ 's candidate distance to be $d+L$.
- Update v's parent to be $u$.
- Update $v$ 's priority in the priority queue to $d+L+h(v, t)$.


## A* Solving Super Mario (video)

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## Minimum Spanning Tree


A. 0-1
D. 6-7
B. 2-3
E. $>7$
C. 4-5

## Prim's Algorithm

## Prim's algorithm

Arbitrarily choose start vertex
Add start vertex to MST
While vertices in MST < total vertices:

- Examine all edges that leave the current MST
- Choose the smallest one
- Add the end vertex of that edge to the MST

Prim's algorithm

## Kruskal's Algorithm

## Kruskal's algorithm

Remove all edges from graph
Place all edges in a PQ based on length/weight
While !PQ.isEmpty():

- Dequeue edge
- If the edge connects previous disconnected nodes or groups of nodes, keep the edge
- Otherwise discard the edge

Kruskal's algorithm

## Kruskal's algorithm

Remove all edges from graph
Place all edges in a PQ based on length/weight While !PQ.isEmpty():

- Dequeue edge

Efficiency of this step is key

- If the edge connects previous disconnected nodes or groups of nodes, keep the edge
- Otherwise discard the edge


## Cluster management questions

The assignment handout asks you to consider questions such as:

- How will you keep track of which nodes are in each cluster?
- How will you determine which cluster a node belongs to?
- How will you merge together two clusters?


## Cluster management strategies

[watch lecture for whiteboard hints]

## The Good Will Hunting Problem

## Video Clip

https://www.youtube.com/watch?v=N7b0cLn-wHU
"Draw all the homeomorphically irreducible trees with $\mathrm{n}=10$."

"Draw all the homeomorphically irreducible trees with $\mathrm{n}=10$."

In this case "trees" simply means graphs with no cycles "with $\mathrm{n}=10$ " (i.e., has 10 nodes)
"homeomorphically irreducible"

- No nodes of degree 2 allowed in your solutions
, For this problem, nodes of degree 2 are useless in terms of tree structure-they just act as a blip on an edge—and are therefore banned
- Have to be actually different
, Ignore superficial changes in rotation or angles of drawing

