# Distributed Optimization via Alternating Direction Method of Multipliers 

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source:
Distributed Optimization and Statistical Learning via the Alternating Direction Method of Multipliers (Boyd, Parikh, Chu, Peleato, Eckstein)

## Goals

robust methods for

- arbitrary-scale optimization
- machine learning/statistics with huge data-sets
- dynamic optimization on large-scale network
- computer vision
- decentralized optimization
- devices/processors/agents coordinate to solve large problem, by passing relatively small messages


## Outline

Dual decomposition
Method of multipliers
Alternating direction method of multipliers
Common patterns

Examples

Consensus and exchange
Conclusions

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## Dual problem

- convex equality constrained optimization problem

$$
\begin{array}{ll}
\operatorname{minimize} & f(x) \\
\text { subject to } & A x=b
\end{array}
$$

- Lagrangian: $L(x, y)=f(x)+y^{T}(A x-b)$
- dual function: $g(y)=\inf _{x} L(x, y)$
- dual problem: maximize $g(y)$
- recover $x^{\star}=\operatorname{argmin}_{x} L\left(x, y^{\star}\right)$


## Dual ascent

- gradient method for dual problem: $y^{k+1}=y^{k}+\alpha^{k} \nabla g\left(y^{k}\right)$
- $\nabla g\left(y^{k}\right)=A \tilde{x}-b$, where $\tilde{x}=\operatorname{argmin}_{x} L\left(x, y^{k}\right)$
- dual ascent method is

$$
\begin{array}{rlr}
x^{k+1} & :=\operatorname{argmin}_{x} L\left(x, y^{k}\right) & / / x \text {-minimization } \\
y^{k+1} & :=y^{k}+\alpha^{k}\left(A x^{k+1}-b\right) & \text { // dual update }
\end{array}
$$

- works, with lots of strong assumptions


## Dual decomposition

- suppose $f$ is separable:

$$
f(x)=f_{1}\left(x_{1}\right)+\cdots+f_{N}\left(x_{N}\right), \quad x=\left(x_{1}, \ldots, x_{N}\right)
$$

- then $L$ is separable in $x: L(x, y)=L_{1}\left(x_{1}, y\right)+\cdots+L_{N}\left(x_{N}, y\right)-y^{T} b$,

$$
L_{i}\left(x_{i}, y\right)=f_{i}\left(x_{i}\right)+y^{T} A_{i} x_{i}
$$

- $x$-minimization in dual ascent splits into $N$ separate minimizations

$$
x_{i}^{k+1}:=\underset{x_{i}}{\operatorname{argmin}} L_{i}\left(x_{i}, y^{k}\right)
$$

which can be carried out in parallel

## Dual decomposition

- dual decomposition (Everett, Dantzig, Wolfe, Benders 1960-65)

$$
\begin{aligned}
x_{i}^{k+1} & :=\operatorname{argmin}_{x_{i}} L_{i}\left(x_{i}, y^{k}\right), \quad i=1, \ldots, N \\
y^{k+1} & :=y^{k}+\alpha^{k}\left(\sum_{i=1}^{N} A_{i} x_{i}^{k+1}-b\right)
\end{aligned}
$$

- scatter $y^{k}$; update $x_{i}$ in parallel; gather $A_{i} x_{i}^{k+1}$
- solve a large problem
- by iteratively solving subproblems (in parallel)
- dual variable update provides coordination
- works, with lots of assumptions; often slow


## Outline

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## Alternating direction method of multipliers

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Method of multipliers

## Method of multipliers

- a method to robustify dual ascent
- use augmented Lagrangian (Hestenes, Powell 1969), $\rho>0$

$$
L_{\rho}(x, y)=f(x)+y^{T}(A x-b)+(\rho / 2)\|A x-b\|_{2}^{2}
$$

- method of multipliers (Hestenes, Powell; analysis in Bertsekas 1982)

$$
\begin{aligned}
x^{k+1} & :=\underset{x}{\operatorname{argmin}} L_{\rho}\left(x, y^{k}\right) \\
y^{k+1} & :=y^{k}+\rho\left(A x^{k+1}-b\right)
\end{aligned}
$$

(note specific dual update step length $\rho$ )

## Method of multipliers dual update step

- optimality conditions (for differentiable $f$ ):

$$
A x^{\star}-b=0, \quad \nabla f\left(x^{\star}\right)+A^{T} y^{\star}=0
$$

(primal and dual feasibility)

- since $x^{k+1}$ minimizes $L_{\rho}\left(x, y^{k}\right)$

$$
\begin{aligned}
0 & =\nabla_{x} L_{\rho}\left(x^{k+1}, y^{k}\right) \\
& =\nabla_{x} f\left(x^{k+1}\right)+A^{T}\left(y^{k}+\rho\left(A x^{k+1}-b\right)\right) \\
& =\nabla_{x} f\left(x^{k+1}\right)+A^{T} y^{k+1}
\end{aligned}
$$

- dual update $y^{k+1}=y^{k}+\rho\left(x^{k+1}-b\right)$ makes $\left(x^{k+1}, y^{k+1}\right)$ dual feasible
- primal feasibility achieved in limit: $A x^{k+1}-b \rightarrow 0$


## Method of multipliers

(compared to dual decomposition)

- good news: converges under much more relaxed conditions ( $f$ can be nondifferentiable, take on value $+\infty, \ldots$ )
- bad news: quadratic penalty destroys splitting of the $x$-update, so can't do decomposition


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## Alternating direction method of multipliers

- a method
- with good robustness of method of multipliers
- which can support decomposition
- "robust dual decomposition" or "decomposable method of multipliers"
- proposed by Gabay, Mercier, Glowinski, Marrocco in 1976


## Alternating direction method of multipliers

- ADMM problem form (with $f, g$ convex)

$$
\begin{array}{ll}
\operatorname{minimize} & f(x)+g(z) \\
\text { subject to } & A x+B z=c
\end{array}
$$

- two sets of variables, with separable objective
- $L_{\rho}(x, z, y)=f(x)+g(z)+y^{T}(A x+B z-c)+(\rho / 2)\|A x+B z-c\|_{2}^{2}$
- ADMM:

$$
\begin{array}{rlr}
x^{k+1} & :=\operatorname{argmin}_{x} L_{\rho}\left(x, z^{k}, y^{k}\right) & / / x \text {-minimization } \\
z^{k+1} & :=\operatorname{argmin}_{z} L_{\rho}\left(x^{k+1}, z, y^{k}\right) & / / z \text {-minimization } \\
y^{k+1} & :=y^{k}+\rho\left(A x^{k+1}+B z^{k+1}-c\right) & / / \text { dual update }
\end{array}
$$

## Alternating direction method of multipliers

- if we minimized over $x$ and $z$ jointly, reduces to method of multipliers
- instead, we do one pass of a Gauss-Seidel method
- we get splitting since we minimize over $x$ with $z$ fixed, and vice versa


## ADMM and optimality conditions

- optimality conditions (for differentiable case):
- primal feasibility: $A x+B z-c=0$
- dual feasibility: $\nabla f(x)+A^{T} y=0, \quad \nabla g(z)+B^{T} y=0$
- since $z^{k+1}$ minimizes $L_{\rho}\left(x^{k+1}, z, y^{k}\right)$ we have

$$
\begin{aligned}
0 & =\nabla g\left(z^{k+1}\right)+B^{T} y^{k}+\rho B^{T}\left(A x^{k+1}+B z^{k+1}-c\right) \\
& =\nabla g\left(z^{k+1}\right)+B^{T} y^{k+1}
\end{aligned}
$$

- so with ADMM dual variable update, $\left(x^{k+1}, z^{k+1}, y^{k+1}\right)$ satisfies second dual feasibility condition
- primal and first dual feasibility are achieved as $k \rightarrow \infty$


## ADMM with scaled dual variables

- combine linear and quadratic terms in augmented Lagrangian

$$
\begin{aligned}
L_{\rho}(x, z, y) & =f(x)+g(z)+y^{T}(A x+B z-c)+(\rho / 2)\|A x+B z-c\|_{2}^{2} \\
& =f(x)+g(z)+(\rho / 2)\|A x+B z-c+u\|_{2}^{2}+\text { const. }
\end{aligned}
$$

with $u^{k}=(1 / \rho) y^{k}$

- ADMM (scaled dual form):

$$
\begin{aligned}
x^{k+1} & :=\underset{x}{\operatorname{argmin}}\left(f(x)+(\rho / 2)\left\|A x+B z^{k}-c+u^{k}\right\|_{2}^{2}\right) \\
z^{k+1} & :=\underset{z}{\operatorname{argmin}}\left(g(z)+(\rho / 2)\left\|A x^{k+1}+B z-c+u^{k}\right\|_{2}^{2}\right) \\
u^{k+1} & :=u^{k}+\left(A x^{k+1}+B z^{k+1}-c\right)
\end{aligned}
$$

## Convergence

- assume (very little!)
- $f, g$ convex, closed, proper
- $L_{0}$ has a saddle point
- then ADMM converges:
- iterates approach feasibility: $A x^{k}+B z^{k}-c \rightarrow 0$
- objective approaches optimal value: $f\left(x^{k}\right)+g\left(z^{k}\right) \rightarrow p^{\star}$


## Related algorithms

- operator splitting methods
(Douglas, Peaceman, Rachford, Lions, Mercier, ... 1950s, 1979)
- Dykstra's alternating projections algorithm (1983)
- Spingarn's method of partial inverses (1985)
- Rockafellar-Wets progressive hedging (1991)
- proximal methods (Rockafellar, many others, 1976-)
- saddle-point proximal methods (Chambolle, Pock 2005-)
- Bregman iterative methods (2008-)
- most of these are special cases of the proximal point algorithm (Rockafellar 1976)


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- $x$-update step requires minimizing $f(x)+(\rho / 2)\|A x-v\|_{2}^{2}$ (with $v=B z^{k}-c+u^{k}$, which is constant during $x$-update)
- similar for $z$-update
- several special cases come up often
- can simplify update by exploiting structure in these cases


## Decomposition

- suppose $f$ is block-separable,

$$
f(x)=f_{1}\left(x_{1}\right)+\cdots+f_{N}\left(x_{N}\right), \quad x=\left(x_{1}, \ldots, x_{N}\right)
$$

- $A$ is conformably block separable: $A^{T} A$ is block diagonal
- then $x$-update splits into $N$ parallel updates of $x_{i}$


## Proximal operator

- consider $x$-update when $A=I$

$$
x^{+}=\underset{x}{\operatorname{argmin}}\left(f(x)+(\rho / 2)\|x-v\|_{2}^{2}\right)=\operatorname{prox}_{f, \rho}(v)
$$

- some special cases:

$$
\begin{array}{cl}
\left.f=I_{C} \text { (indicator fct. of set } C\right) & x^{+}:=\Pi_{C}(v) \text { (projection onto } C \text { ) } \\
f=\lambda\|\cdot\|_{1}\left(\ell_{1}\right. \text { norm) } & x_{i}^{+}:=S_{\lambda / \rho}\left(v_{i}\right) \text { (soft thresholding) } \\
\left(S_{a}(v)=(v-a)_{+}-(-v-a)_{+}\right) &
\end{array}
$$

## Quadratic objective

- $f(x)=(1 / 2) x^{T} P x+q^{T} x+r$
- $x^{+}:=\left(P+\rho A^{T} A\right)^{-1}\left(\rho A^{T} v-q\right)$
- use matrix inversion lemma when computationally advantageous

$$
\left(P+\rho A^{T} A\right)^{-1}=P^{-1}-\rho P^{-1} A^{T}\left(I+\rho A P^{-1} A^{T}\right)^{-1} A P^{-1}
$$

- (direct method) cache factorization of $P+\rho A^{T} A$ (or $\left.I+\rho A P^{-1} A^{T}\right)$
- (iterative method) warm start, early stopping, reducing tolerances


## Smooth objective

- $f$ smooth
- can use standard methods for smooth minimization
- gradient, Newton, or quasi-Newton
- preconditionned CG, limited-memory BFGS (scale to very large problems)
- can exploit
- warm start
- early stopping, with tolerances decreasing as ADMM proceeds


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## Constrained convex optimization

- consider ADMM for generic problem

$$
\begin{array}{ll}
\operatorname{minimize} & f(x) \\
\text { subject to } & x \in \mathcal{C}
\end{array}
$$

- ADMM form: take $g$ to be indicator of $\mathcal{C}$

$$
\begin{array}{ll}
\operatorname{minimize} & f(x)+g(z) \\
\text { subject to } & x-z=0
\end{array}
$$

- algorithm:

$$
\begin{aligned}
x^{k+1} & :=\underset{x}{\operatorname{argmin}}\left(f(x)+(\rho / 2)\left\|x-z^{k}+u^{k}\right\|_{2}^{2}\right) \\
z^{k+1} & :=\Pi_{\mathcal{C}}\left(x^{k+1}+u^{k}\right) \\
u^{k+1} & :=u^{k}+x^{k+1}-z^{k+1}
\end{aligned}
$$

## Lasso

- lasso problem:

$$
\operatorname{minimize} \quad(1 / 2)\|A x-b\|_{2}^{2}+\lambda\|x\|_{1}
$$

- ADMM form:

$$
\begin{array}{ll}
\operatorname{minimize} & (1 / 2)\|A x-b\|_{2}^{2}+\lambda\|z\|_{1} \\
\text { subject to } & x-z=0
\end{array}
$$

- ADMM:

$$
\begin{aligned}
x^{k+1} & :=\left(A^{T} A+\rho I\right)^{-1}\left(A^{T} b+\rho z^{k}-y^{k}\right) \\
z^{k+1} & :=S_{\lambda / \rho}\left(x^{k+1}+y^{k} / \rho\right) \\
y^{k+1} & :=y^{k}+\rho\left(x^{k+1}-z^{k+1}\right)
\end{aligned}
$$

## Lasso example

- example with dense $A \in \mathbf{R}^{1500 \times 5000}$ (1500 measurements; 5000 regressors)
- computation times

| factorization (same as ridge regression) | 1.3 s |
| :--- | :--- |
| subsequent ADMM iterations | 0.03 s |
| lasso solve (about 50 ADMM iterations) | 2.9 s |
| full regularization path $(30 \lambda$ 's $)$ | 4.4 s |

- not bad for a very short Matlab script


## Sparse inverse covariance selection

- S: empirical covariance of samples from $\mathcal{N}(0, \Sigma)$, with $\Sigma^{-1}$ sparse (i.e., Gaussian Markov random field)
- estimate $\Sigma^{-1}$ via $\ell_{1}$ regularized maximum likelihood minimize $\operatorname{Tr}(S X)-\log \operatorname{det} X+\lambda\|X\|_{1}$
- methods: COVSEL (Banerjee et al 2008), graphical lasso (FHT 2008)


## Sparse inverse covariance selection via ADMM

- ADMM form:

$$
\begin{array}{ll}
\operatorname{minimize} & \operatorname{Tr}(S X)-\log \operatorname{det} X+\lambda\|Z\|_{1} \\
\text { subject to } & X-Z=0
\end{array}
$$

- ADMM:

$$
\begin{aligned}
X^{k+1} & :=\underset{X}{\operatorname{argmin}}\left(\operatorname{Tr}(S X)-\log \operatorname{det} X+(\rho / 2)\left\|X-Z^{k}+U^{k}\right\|_{F}^{2}\right) \\
Z^{k+1} & :=S_{\lambda / \rho}\left(X^{k+1}+U^{k}\right) \\
U^{k+1} & :=U^{k}+\left(X^{k+1}-Z^{k+1}\right)
\end{aligned}
$$

## Analytical solution for $X$-update

- compute eigendecomposition $\rho\left(Z^{k}-U^{k}\right)-S=Q \Lambda Q^{T}$
- form diagonal matrix $\tilde{X}$ with

$$
\tilde{X}_{i i}=\frac{\lambda_{i}+\sqrt{\lambda_{i}^{2}+4 \rho}}{2 \rho}
$$

- let $X^{k+1}:=Q \tilde{X} Q^{T}$
- cost of $X$-update is an eigendecomposition


## Sparse inverse covariance selection example

- $\Sigma^{-1}$ is $1000 \times 1000$ with $10^{4}$ nonzeros
- graphical lasso (Fortran): 20 seconds -3 minutes
- ADMM (Matlab): 3-10 minutes
- (depends on choice of $\lambda$ )
- very rough experiment, but with no special tuning, ADMM is in ballpark of recent specialized methods
- (for comparison, COVSEL takes $25+\min$ when $\Sigma^{-1}$ is a $400 \times 400$ tridiagonal matrix)


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## Consensus optimization

- want to solve problem with $N$ objective terms

$$
\operatorname{minimize} \quad \sum_{i=1}^{N} f_{i}(x)
$$

- e.g., $f_{i}$ is the loss function for $i$ th block of training data
- ADMM form:

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i=1}^{N} f_{i}\left(x_{i}\right) \\
\text { subject to } & x_{i}-z=0
\end{array}
$$

- $x_{i}$ are local variables
- $z$ is the global variable
- $x_{i}-z=0$ are consistency or consensus constraints
- can add regularization using a $g(z)$ term


## Consensus optimization via ADMM

- $L_{\rho}(x, z, y)=\sum_{i=1}^{N}\left(f_{i}\left(x_{i}\right)+y_{i}^{T}\left(x_{i}-z\right)+(\rho / 2)\left\|x_{i}-z\right\|_{2}^{2}\right)$
- ADMM:

$$
\begin{aligned}
x_{i}^{k+1} & :=\underset{x_{i}}{\operatorname{argmin}}\left(f_{i}\left(x_{i}\right)+y_{i}^{k T}\left(x_{i}-z^{k}\right)+(\rho / 2)\left\|x_{i}-z^{k}\right\|_{2}^{2}\right) \\
z^{k+1} & :=\frac{1}{N} \sum_{i=1}^{N}\left(x_{i}^{k+1}+(1 / \rho) y_{i}^{k}\right) \\
y_{i}^{k+1} & :=y_{i}^{k}+\rho\left(x_{i}^{k+1}-z^{k+1}\right)
\end{aligned}
$$

- with regularization, averaging in $z$ update is followed by $\operatorname{prox}_{g, \rho}$


## Consensus optimization via ADMM

- using $\sum_{i=1}^{N} y_{i}^{k}=0$, algorithm simplifies to

$$
\begin{aligned}
& x_{i}^{k+1}:=\underset{x_{i}}{\operatorname{argmin}}\left(f_{i}\left(x_{i}\right)+y_{i}^{k T} x_{i}+(\rho / 2)\left\|x_{i}-\bar{x}^{k}\right\|_{2}^{2}\right) \\
& y_{i}^{k+1}:=y_{i}^{k}+\rho\left(x_{i}^{k+1}-\bar{x}^{k+1}\right) \\
& \text { where } \bar{x}^{k}=(1 / N) \sum_{i=1}^{N} x_{i}^{k}
\end{aligned}
$$

- in each iteration
- gather $x_{i}^{k}$ and average to get $\bar{x}^{k}$
- scatter the average $\bar{x}^{k}$ to processors
- update $y_{i}^{k}$ locally (in each processor, in parallel)
- update $x_{i}$ locally


## Statistical interpretation

- $f_{i}$ is negative log-likelihood for parameter $x$ given $i$ th data block
- $x_{i}^{k+1}$ is MAP estimate under prior $\mathcal{N}\left(\bar{x}^{k}+(1 / \rho) y_{i}^{k}, \rho I\right)$
- prior mean is previous iteration's consensus shifted by 'price' of processor $i$ disagreeing with previous consensus
- processors only need to support a Gaussian MAP method
- type or number of data in each block not relevant
- consensus protocol yields global maximum-likelihood estimate


## Consensus classification

- data (examples) $\left(a_{i}, b_{i}\right), i=1, \ldots, N, a_{i} \in \mathbf{R}^{n}, b_{i} \in\{-1,+1\}$
- linear classifier $\operatorname{sign}\left(a^{T} w+v\right)$, with weight $w$, offset $v$
- margin for $i$ th example is $b_{i}\left(a_{i}^{T} w+v\right)$; want margin to be positive
- loss for $i$ th example is $l\left(b_{i}\left(a_{i}^{T} w+v\right)\right)$
- $l$ is loss function (hinge, logistic, probit, exponential, ...)
- choose $w, v$ to minimize $\frac{1}{N} \sum_{i=1}^{N} l\left(b_{i}\left(a_{i}^{T} w+v\right)\right)+r(w)$
- $r(w)$ is regularization term $\left(\ell_{2}, \ell_{1}, \ldots\right)$
- split data and use ADMM consensus to solve


## Consensus SVM example

- hinge loss $l(u)=(1-u)_{+}$with $\ell_{2}$ regularization
- baby problem with $n=2, N=400$ to illustrate
- examples split into 20 groups, in worst possible way: each group contains only positive or negative examples

Iteration 1


Consensus and exchange

Iteration 5


## Iteration 40



## Distributed lasso example

- example with dense $A \in \mathbf{R}^{400000 \times 8000}$ (roughly 30 GB of data)
- distributed solver written in C using MPI and GSL
- no optimization or tuned libraries (like ATLAS, MKL)
- split into 80 subsystems across 10 (8-core) machines on Amazon EC2
- computation times

| loading data | 30 s |
| :--- | :--- |
| factorization | 5 m |
| subsequent ADMM iterations | $0.5-2 \mathrm{~s}$ |
| lasso solve (about 15 ADMM iterations) | $5-6 \mathrm{~m}$ |

## Exchange problem

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i=1}^{N} f_{i}\left(x_{i}\right) \\
\text { subject to } & \sum_{i=1}^{N} x_{i}=0
\end{array}
$$

- another canonical problem, like consensus
- in fact, it's the dual of consensus
- can interpret as $N$ agents exchanging $n$ goods to minimize a total cost
- $\left(x_{i}\right)_{j} \geq 0$ means agent $i$ receives $\left(x_{i}\right)_{j}$ of good $j$ from exchange
- $\left(x_{i}\right)_{j}<0$ means agent $i$ contributes $\left|\left(x_{i}\right)_{j}\right|$ of good $j$ to exchange
- constraint $\sum_{i=1}^{N} x_{i}=0$ is equilibrium or market clearing constraint
- optimal dual variable $y^{\star}$ is a set of valid prices for the goods


## Exchange ADMM

- solve as a generic constrained convex problem with constraint set

$$
\mathcal{C}=\left\{x \in \mathbf{R}^{n N} \mid x_{1}+x_{2}+\cdots+x_{N}=0\right\}
$$

- scaled form:

$$
\begin{aligned}
x_{i}^{k+1} & :=\underset{x_{i}}{\operatorname{argmin}}\left(f_{i}\left(x_{i}\right)+(\rho / 2)\left\|x_{i}-x_{i}^{k}+\bar{x}^{k}+u^{k}\right\|_{2}^{2}\right) \\
u^{k+1} & :=u^{k}+\bar{x}^{k+1}
\end{aligned}
$$

- unscaled form:

$$
\begin{aligned}
x_{i}^{k+1} & :=\underset{x_{i}}{\operatorname{argmin}}\left(f_{i}\left(x_{i}\right)+y^{k T} x_{i}+(\rho / 2)\left\|x_{i}-\left(x_{i}^{k}-\bar{x}^{k}\right)\right\|_{2}^{2}\right) \\
y^{k+1} & :=y^{k}+\rho \bar{x}^{k+1}
\end{aligned}
$$

## Interpretation as tâtonnement process

- tâtonnement process: iteratively update prices to clear market
- work towards equilibrium by increasing/decreasing prices of goods based on excess demand/supply
- dual decomposition is the simplest tâtonnement algorithm
- ADMM adds proximal regularization
- incorporate agents' prior commitment to help clear market
- convergence far more robust convergence than dual decomposition


## Distributed dynamic energy management

- $N$ devices exchange power in time periods $t=1, \ldots, T$
- $x_{i} \in \mathbf{R}^{T}$ is power flow profile for device $i$
- $f_{i}\left(x_{i}\right)$ is cost of profile $x_{i}$ (and encodes constraints)
- $x_{1}+\cdots+x_{N}=0$ is energy balance (in each time period)
- dynamic energy management problem is exchange problem
- exchange ADMM gives distributed method for dynamic energy management
- each device optimizes its own profile, with quadratic regularization for coordination
- residual (energy imbalance) is driven to zero


## Example

- network with 8000 devices exchanging power at 3000 nodes (mixture of generators, batteries, smart loads, transmission lines, ...)
- coordinate devices over 96 time periods
- $\sim 1$ million variables in optimization problem



Consensus and exchange

## Solve time scaling



- serial multi-threaded implementation on 32-core machine with 64 independent threads
- best fit exponent is 0.996
- fully decentralized computation would result in sub second solve time for any size network


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## Summary and conclusions

ADMM

- is the same as, or closely related to, many methods with other names
- has been around since the 1970s
- gives simple single-processor algorithms that can be competitive with state-of-the-art
- can be used to coordinate many processors, each solving a substantial problem, to solve a very large problem

